

Load Flow as a Planning Tool - An Optimal Approach

By

R. Shanthi

M. Jayakumar

V. S. Ravichandran

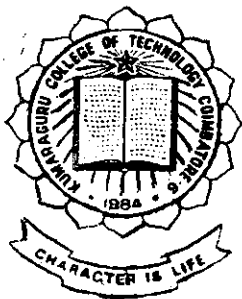
G. Shanthanayaki

Shiny Purushothaman

P-28

Guided By

Er K. Surendran



Project Work

1989

Faculty of Electrical & Electronics Engineering
Kumaraguru College of Technology

Coimbatore - 641 006

Faculty of Electrical and Electronics Engineering
Kumaraguru College of Technology

Coimbatore-641 006

Certificate

This is the Bonafide Record of the Project titled
Load Flow as a Planning Tool - An Optimal Approach

done by

Mr. / Miss. _____ Roll No. _____

In partial fulfilment of the requirements for the Degree of
Bachelor of Engineering in Electrical and Electronics Engineering
Branch of Bharathiar University, Coimbatore-641 046
during the year 1988-89

Dr. K. A. PALANISWAMY, B.E., M.Sc. (Engg.), Ph.D.,
MISTE, C. Engg (I), FIEE,

Professor and Head

Department of Electrical and Electronics Engineering,
Kumaraguru College of Technology,
(Head of Department) 006

Er. K. Surendran
(Project Guide)

Submitted for the University Examination held on _____

University Register No. _____

Internal Examiner

External Examiner

*Dedicated
To
Our Parents*

acknowledgement

ACKNOWLEDGEMENT

At the juncture of accomplishment of our project, we acknowledge with due courtesy **Mr. K. SURENDRAN M.Tech.**, for his guidance and inspiration. We are grateful to our Professor **Mr.K.A. PALANISWAMY Ph. D, M.I.S.T.E., F.I.E.**, for his valuable suggestions and encouragement.

It is our privelege to thank Principal Professor, **R.M. LAKSHMANAN, B.E., M.Sc., (Engg), M.I.S.T.E., M.S.A., E.S.T.**, for his co-operation in our endeavour.

Our sincere thanks are due to Professor **P.SHANMUGAM, M.Sc., (Engg), M.S. (HOWAI) M.I.E.E.**, for the help rendered to us.

We are greatly indebted to **Mr. K.RAMAN NAIR B.Sc., (Engg) M.Tech., Ph.D.**, for suggesting us this project and for his constant help and guidance.

We extend our deep sense of gratitude to **Mr. R.P. THIANGARAJ M.Sc., M.Phil.**, who helped us a lot in drafting the software techniques of our project.

A special word of thanks is, due to **Miss. MUTHAMANI M.E.**, for her encouragement throughout our project.

A final work of thanks is extended to the faculty members of Electrical and Electronics Department for their whole hearted co-operation in bringing this project into shape.

abstract

LOAD FLOW AS A PLANNING TOOL.- AN OPTIMAL APPROACH

ABSTRACT:

The load flow solution is an indispensable tool in the design or operational planning of a power system. In practical systems, in order to obtain an acceptable design or operating strategy, a large number of conventional load flow solutions have to be performed one after the other. The success of the final solution depends to a great extent on the experience, skill and intuition of the system planner and invariably involves trial and error adjustments.

In this Report, a new approach is conceived wherein a single optimal load flow solution replaces a series of many trial and error load flow solutions. For a system under given constraints the best possible operating strategy is directly found by using one line solution technique. This method is illustrated with an example. While the proposed technique calls for greater programming effort and computer memory compared to conventional load flows, this would not be a serious drawback with the availability of the present day powerful computing aids.

nomenclature

NOMENCLATURE

P	- Real Power
Q	- Reactive Power
V	- Voltage magnitude
	- Voltage angle in radians
N	- Number of buses
G	- Number of generator buses

SUPER SCRIPT

-1	- Inverse of the Matrix
tr	- Transpose of the Matrix

contents

introduction

CONTENTS

CHAPTER		PAGE NO
	INTRODUCTION	1
	GENERAL	5
1	OPTIMIZATION TECHNIQUES	7
2	WEIGHTED LEAST SQUARE ERROR ALGORITHM	9
3	PROBLEM FORMULATION	13
4	DERIVATION OF THE ALGORITHM	27
5	ITERATIVE PROCEDURE	31
6	FLOW DIAGRAM	49
7	SOFTWARE AND RESULTS	65
8	FORMULA USED	121
9	CONCLUSION	129
10	BIBLIOGRAPHY	131

INTRODUCTION

The load flow solution techniques based on Gauss Seidal or Newton Raphson Algorithms are so frequently used in the process of power system design or operational planning so that its basic features and limitations have been taken for granted. Each bus in a system is characterised by four variables, $|V|$ and δ at the slack bus, P and $|V|$ at generator buses and P and Q at the load buses. The load flow solution gives the corresponding values of P , and Q at the slack bus, Q and δ at the generator buses and $|V|$ and δ at the load buses. Often the results of the initial load flow are not acceptable. For example the P and Q found at the slack bus may be beyond the capacity of the generator at that bus. Similarly $|V|$ as found at some of the load buses may be too low. Or at a generator bus, due to the violation of limits on Q , the voltage may be too high or too low. By reducing P at such a bus, the limits for Q may be widened and this could lead to an acceptable voltage level. Thus, the system planner has to consider several alterations, modify the input specifications to the load flow problem and obtain a fresh solution. Often a number of such load flows may be carried out before an acceptable result is obtained.

In this work, a new method is proposed which enable the system planner to specify desired values to any number of

general

solution.

By this method the optimum results are achieved with a

optimization techniques

variables in the system. The load flow solution is reconceived as one that yields a system state that is closest to the desired solution. To solve the problem, an optimization technique based on weighted least square errors principle is used. Provision are also made for the system planner to incorporate the relative importance that he attaches to each of the various desired values. The method yields a solution in which the sum of the weighted squares of the deviation of each variable from its desired or target value is minimum subject to the necessary load flow constraints. Without the need to conduct several case studies. Thus the process of arriving at the most acceptable system state is made more systematic and objective. And less time consuming. and the subjective decisions which the system planner normally has to make are almost entirely eliminated.

CHAPTER I

OPTIMIZATION TECHNIQUES:

Art or technique of obtaining best results under certain restrictions and limits is known as the optimization. The objective function may be different from problem to problem. The method for getting the results may also vary. In general the optimization techniques fall into two categories, viz Linear and Nonlinear techniques. Out of these two, linear methods are more effective, less time consuming and for real time implementation they are well suited. In this chapter weighted least square error method is reviewed briefly. A feasible solution which optimizes the objective function is called an optimal solution.

weighted least square error algorithm

This result is also in accord with experience and common sense.

If the accidental errors of n measurements m_1, m_2, \dots, m_n be denoted by x_1, x_2, \dots, x_n respectively then

$$x_1^2 + x_2^2 + x_3^2 + \dots + x_n^2 = \sum x^2 \text{ is a minimum.}$$

We arrive at this principle from the principle of probability, which states that the best or most probable value obtainable from a set of measurements or observations of equal precision is that value for which the sum of the squares of the errors is a minimum.

In its most general form, the above principle states that the best value of an unknown quantity that can be obtained from a set of measurements of unequal precision is that which makes the sum of the weighted squares of the errors a minimum.

Adopting this principle to the optimal load flow problem, we define the best state of the power system as that for which the sum of the weighted squares of the deviations of each variable from its derived value is a minimum. The weightages allow the system planner to assign the relative importance of the various variables. This is in fact the only point at which the subjective assessment of the system planner can affect the solution.

problem formulation

CHAPTER III

PROBLEM FORMULATION:

Problem is formulated in polar co-ordinates. Variables at each bus are P , Q , V and δ

$N \rightarrow$ total number of buses

1 to $G \rightarrow$ Generator buses of which bus 1 is slack bus, so that $S_1 = 0$

$G + 1$ to $N \rightarrow$ load buses, at which P and Q are specified.

δ_i , $i = 2$ to N and V_i , $i = 1$ to N are the system state variables numbering $2N - 1$. There are $2N$ other variables, namely P_i , $i = 1$ to N and Q_i , $i = 1$ to N .

In conventional load flow studies one has to specify values for $2N$ of the $4N$ variables and then solve the non linear power equations for the remaining $2N$ variables. However in practice we really need to specify only the slack bus angle and P and Q for the load buses. All other specifications are artificially introduced.

Let there be G generator buses, including the slack bus. Then $G+1, G+2, \dots, N$ are the load buses. For each of these buses, P_i and Q_i must be specified. So the total number of specified variables is $K = 2(N-G)$.

From among the $2N-1$ state variables, we choose K variables called dependent variables, forming a vector $[x]$ of order $K \times 1$.

$$\begin{array}{l}
 1 \\
 2 \\
 \vdots \\
 N-G \\
 N-G+1 \\
 \vdots \\
 \vdots \\
 K=2(N-G)
 \end{array}
 \begin{array}{c}
 \delta_{G+1} \\
 \delta_{G+1} \\
 \vdots \\
 \vdots \\
 \delta_N \\
 V_{G+1} \\
 V_{G+2} \\
 \vdots \\
 \vdots \\
 V_N
 \end{array}$$

The remaining $2N-1-k$ variables is δ_2 to δ_G , V_1 to V_G are called independent variables and form a vector $[Y]$ of order $j \times 1$ where $j = 2N-1-k = 2G-1$

$$\begin{array}{l}
 1 \\
 2 \\
 \vdots \\
 \vdots \\
 \vdots \\
 G-1 \\
 G \\
 G \\
 \vdots \\
 \vdots \\
 j=2G-1
 \end{array}
 \begin{array}{c}
 \delta_2 \\
 \delta_3 \\
 \vdots \\
 \vdots \\
 \vdots \\
 \delta_G \\
 \text{---} \\
 V_1 \\
 V_2 \\
 \vdots \\
 \vdots \\
 V_G
 \end{array}$$

$$[Q] = \begin{bmatrix} Q_1 \\ Q_2 \\ \cdot \\ \cdot \\ Q_G \end{bmatrix} \quad [Q_{\min}] = \begin{bmatrix} Q_{1\min} \\ Q_{2\min} \\ \cdot \\ \cdot \\ Q_{G\min} \end{bmatrix} \quad [Q_{\max}] = \begin{bmatrix} Q_{1\max} \\ Q_{2\max} \\ \cdot \\ \cdot \\ Q_{G\max} \end{bmatrix}$$

There is a third class of variables for which we want to have values as close as possible to give 'target' values while also lying between specified lower and upper limits. This includes P at each generator bus and V for all buses. These 'target' variables are included in a vector [b] of order m X 1 where m = N+G. Their target values are included in a vector [t] and their limits in vectors [b_{min}] and [b_{max}] all of order m x 1.

$$[b] = \begin{bmatrix} 1 & P_1 \\ 2 & P_2 \\ \cdot & \cdot \\ \cdot & \cdot \\ G & P_G \\ \hline G+1 & V_1 \\ G+2 & V_2 \\ \cdot & \cdot \\ \cdot & \cdot \\ 2G & V_G \\ 2G+1 & V_{G+1} \\ \cdot & \cdot \\ \cdot & \cdot \\ m & V_N \end{bmatrix} \quad [b_{\min}] = \begin{bmatrix} 1 & P_{1\min} \\ 2 & P_{2\min} \\ \cdot & \cdot \\ \cdot & \cdot \\ G & P_{G\min} \\ \hline G+1 & V_{1\min} \\ G+2 & V_{2\min} \\ \cdot & \cdot \\ \cdot & \cdot \\ 2G & V_{G\min} \\ 2G+1 & V_{G+1\min} \\ \cdot & \cdot \\ \cdot & \cdot \\ m & V_{N\min} \end{bmatrix}$$

$$[b_{\max}] =$$

$$\begin{bmatrix} P_{1\max} \\ P_{2\max} \\ \cdot \\ \cdot \\ P_{G\max} \\ V_{1\max} \\ V_{2\max} \\ \cdot \\ \cdot \\ V_{G\max} \\ V_{G+1\max} \\ \cdot \\ \cdot \\ V_{N\max} \end{bmatrix}$$

$$[t] =$$

$$\begin{bmatrix} t_1 \\ t_2 \\ \cdot \\ \cdot \\ t_G \\ t_{G+1} \\ t_{G+2} \\ \cdot \\ \cdot \\ t_{2G} \\ t_{2G+1} \\ \cdot \\ \cdot \\ t_m \end{bmatrix}$$

All the specified variables are included in a vector [a] of order $k \times 1$. The specified values for these variables are grouped in a vector [s] of order $k \times 1$.

$$[a] = \begin{matrix} 1 & P_{G+1} \\ 2 & P_{G+2} \\ \vdots & \vdots \\ N-G & P_N \\ N-G+1 & Q_{G+1} \\ \cdot & Q_{G+2} \\ \cdot & \vdots \\ \cdot & \vdots \\ K & Q_N \end{matrix} \quad [S] = \begin{matrix} S_1 \\ S_2 \\ \vdots \\ S_{N-G} \\ S_{N-G+1} \\ \vdots \\ \vdots \\ S_K \end{matrix}$$

Each element in [a] is a function of the variables [x] and [y]. We may therefore write

$$[a] = [a \{ (x), (y) \}]$$

The load flow solution must satisfy the conditions

$$[a \{ (x), (y) \}] = [s] \quad \text{-----(1)}$$

Now the Q at each generator bus is required to lie between lower and upper limits. The load flow solution must satisfy the conditions.

$$[Q_{\min}] \leq [Q \{ (x), (y) \}] \leq [Q_{\max}] \quad \text{-----(2)}$$

where [Q], [Q_{min}] and [Q_{max}] are $G \times 1$ vectors given by

The load flow solution must satisfy the condition

$$|b_{\min}| \leq |b\{x, y\}| \leq |b_{\max}| \quad (3)$$

and

$$|b\{x, y\}| \text{ must lie as close as possible to } |t| \quad \text{-----(4)}$$

The heuristic condition of equation (4) may be written as a formal mathematical weighted least squares proportion, namely

$$F = \sum_{i=1}^m W_i (t_i - b_i)^2 \text{ must be minimum.}$$

$i=1$

$F = ([t] - [b]^{tr} [W]) ([t] - [b])$ must be minimum where $[W]$ is a $m \times m$ diagonal matrix whose diagonal elements are the weightages w_i

$$[W] = \begin{bmatrix} w_1 & 0 & \dots & 0 \\ 0 & w_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & & w_m \end{bmatrix}$$

The load flow problem may now be formulated to an optimisation problem, as given below. Minimise

$$F = ([t] - [b\{x, y\}]^{tr} [W]) ([t] - [b\{x, y\}]) \quad \text{--- (5)}$$

Subject to the constraints

$$[a\{x, y\}] = [s] \quad \text{-----(1)}$$

$$|Q_{\min}| \leq |Q\{x, y\}| \leq |Q_{\max}| \quad \text{-----(2)}$$

$$|b_{\min}| \leq |b\{x, y\}| \leq |b_{\max}| \quad \text{-----(3)}$$

Note that $K = 2(N-G) = 2N - 2G$

and so $K < 2N - 1$ or $K < \text{no of state variables}$.

Also, $m = N+G = j + N + G + 1$ and so $m > j$ or $m > \text{number of independent variables}$. Since the number of specification K is less than the number of state variables $2N-1$, an infinite number of solution exist which will satisfy equation (1). Out of these, we are looking for that solution which will minimise F of equation (5) and also satisfy and also satisfy the constraints of equation (2) and (3).

derivation of the algorithm

$$= ([t] - [b] - [c] [\Delta x] - [D] [\Delta y])^{\text{tr}} [W] ([t] - [b] - [c] [\Delta x] - [D] [\Delta y])$$

where $[c] = [db/dx]$ is a $m \times k$ matrix

and $[D] = [db/dy]$ is a $m \times j$ matrix

$$= ([t] - [b] - [c] [A]^{-1} ([s] - [a]) - ([D] - [c] [A]^{-1} [B] [\Delta Y])^{\text{tr}} [W] ([t] - [b] - [c] [A]^{-1} ([s] - [a]) - ([D] - [c] [A]^{-1} [B]) [\Delta y])$$

$$= ([t] - [b] - [c] [A]^{-1} ([s] - [a]) - [G] [\Delta y])^{\text{tr}} [W] ([t] - [b] - [c] [A]^{-1} ([s] - [a]) - [G] [\Delta y])$$

where $[G] = [D] - [C] [A]^{-1} [B]$ is a $m \times j$ matrix.

$$= ([t] - [b] - [c] [A]^{-1} ([s] - [a]))^{\text{tr}} [W] ([t] - [b] - [c] [A]^{-1} ([s] - [a])) - 2 ([t] - [b] - [c] [A]^{-1} ([s] - [a]))^{\text{tr}} [W] [G] [\Delta y] + ([G] [\Delta y])^{\text{tr}} [W] ([G] [\Delta y])$$

must be a minimum

therefore $dF/d\Delta Y$ must be zero ie

$$-2 ([t] - [b] - [c] [A]^{-1} ([s] - [a]))^{\text{tr}} [W] [G] + 2 ([G] [\Delta y])^{\text{tr}} [W] [G] = 0$$

transposing

$$-2 [G]^{\text{tr}} [W] ([t] - [b] - [c] [A]^{-1} ([s] - [a])) + 2 [G]^{\text{tr}} [W] ([G] [\Delta y]) = 0$$

ie

$$[G]^{\text{tr}} [W] [G] [\Delta Y] = [G]^{\text{tr}} [W] ([t] - [b] - [c] [A]^{-1} ([s] - [a]))$$

or

$$[\Delta y] = ([G]^{\text{tr}} [W] [G])^{-1} [G]^{\text{tr}} [W] ([t] - [b] - [c] [A]^{-1} ([s] - [a]))$$

CHAPTER IVDERIVATION OF THE ALGORITHM:

Since the load flow problem is non-linear an iterative procedure must be used. At any iteration, let $[x]$ and $[y]$ represent the values of the state variables. Let $[\Delta x]$, $[\Delta y]$ be the optimal corrections to be made to $[x]$ and $[y]$ to proceed to the next iteration.

Then since we are aiming to satisfy equation 1,

$$[S] - [a\{[x] + [\Delta x], [y] + [\Delta y]\}] = 0$$

$$\text{Let } [a\{[x] + [\Delta x], [y] + [\Delta y]\}] = [a\{[x], [y]\}] + [\Delta a]$$

$$\text{Then } [s] - ([a\{[x], [y]\}] + [\Delta a]) = 0$$

$$\text{ie } [S] - [a] - [\Delta a] = 0$$

$$\text{ie } [s] - [a] - [A]\Delta x - [B]\Delta y = 0$$

where $[A] = [da/dx]$ is a $k \times k$ matrix

and $[B] = [da/dy]$ is a $k \times j$ matrix

$$\text{ie } [A][\Delta x] = [s] - [a] - [B][\Delta y]$$

$$\text{pr } [\Delta x] = [A]^{-1} ([s] - [a] - [B][\Delta y]) \quad \text{-----(6)}$$

Also since we are assuming to minimise F ,

from equation 5,

$$F = ([t] - [b\{[x] + [\Delta x], [y] + [\Delta y]\}]) = [b\{[x], [y]\}] + [\Delta b]$$

must be a minimum

$$\text{let } [b\{[x] + [\Delta x], [y] + [\Delta y]\}] = [b\{[x], [y]\}] + [\Delta b]$$

$$\text{then } F = ([t] - [b\{[x], [y]\}] - [\Delta b])^{tr} [W] ([t] - [b\{[x], [y]\}] - [\Delta b])$$

iterative procedure

CHAPTER VITERATIVE PROCEDURE:

The starting values for the state variables [x] and [y] may be chosen as follows:

$$[x]_0 = \begin{bmatrix} \delta_{G+1} \\ \delta_{G+2} \\ \vdots \\ \delta_N \\ \hline V_{G+1} \\ V_{G+2} \\ \vdots \\ V_N \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ \hline 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix}; [y]_0 = \begin{bmatrix} \delta_2 \\ \delta_3 \\ \vdots \\ \delta_G \\ \hline V_1 \\ V_2 \\ \vdots \\ V_G \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ \hline 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix}$$

Let the state variables at the end of an iteration be [x], [y].

To proceed to the next iteration:-

Step 1

Calculate P and Q for all the buses and form the vectors [a], [Q] and [b].

Step 2

The MVA rating of each generator bus is known, say S_i for the i th bus. Calculate $Q_{\text{mini}} = \sqrt{S_i^2 - P_i^2}$ and $Q_{\text{maxi}} = \sqrt{S_i^2 - P_i^2}$

$$\begin{array}{l}
 [a'] = \\
 K' \times 1
 \end{array}
 \begin{array}{c}
 P_{G+1} \\
 P_{G+2} \\
 \vdots \\
 P_N \\
 \text{-----} \\
 Q_{G+1} \\
 Q_{G+2} \\
 \vdots \\
 Q_N \\
 \text{-----} \\
 Q_2 \\
 Q_5
 \end{array}
 \qquad
 \begin{array}{l}
 [s'] = \\
 K' \times 1
 \end{array}
 \begin{array}{c}
 S_1 \\
 S_2 \\
 \vdots \\
 S_{N-G} \\
 \text{-----} \\
 S_{N-G+1} \\
 S_{N-G+2} \\
 \vdots \\
 S_k \\
 \text{-----} \\
 Q_{2min} \\
 Q_{5.max}
 \end{array}$$

If the Q' 's for all the generator buses are within this limits, simply set $K' = K$, $j' = j$, $[x'] = [x]$, $[y'] = [y]$, $[a'] = [a]$ and $[s'] = [s]$.

Step 3

Check each variable in $[b]$ for limit violation. If it is found that $b_i < b_{\text{mini}}$ or $b_i > b_{\text{maxi}}$ increase w_i suitably and thus form a new $[W]$

Step 4

Compute $[s'] - [a']$ and store it in the same location as $[a']$, compute $[t] - [b]$ and store it in the location as $[b]$, compute the $K' \times K'$ matrix $[A']$ as follows.

[A'] =

$\frac{\delta P_{G+1}}{\partial \delta_{G+1}}$	$\frac{\delta P_{G+1}}{\partial \delta_{G+2}}$	\dots	$\frac{\delta P_{G+1}}{\partial \delta_N}$	$\frac{\delta P_{G+1}}{\partial V_{G+1}}$	$\frac{\delta P_{G+1}}{\partial V_{G+2}}$	\dots	$\frac{\delta P_{G+1}}{\partial V_N}$	$\frac{\delta P_{G+1}}{\partial V_2}$	$\frac{\delta P_{G+1}}{\partial V_5}$	1
$\frac{\partial P_{G+2}}{\partial \delta_{G+1}}$	$\frac{\partial P_{G+2}}{\partial \delta_{G+2}}$	\dots	$\frac{\partial P_{G+2}}{\partial \delta_N}$	$\frac{\partial P_{G+2}}{\partial V_{G+1}}$	$\frac{\partial P_{G+2}}{\partial V_{G+2}}$	\dots	$\frac{\partial P_{G+2}}{\partial V_N}$	$\frac{\partial P_{G+2}}{\partial V_2}$	$\frac{\partial P_{G+2}}{\partial V_5}$	2
\vdots										
$\frac{\partial P_N}{\partial \delta_{G+1}}$	$\frac{\partial P_N}{\partial \delta_{G+2}}$	\dots	$\frac{\partial P_N}{\partial \delta_N}$	$\frac{\partial P_N}{\partial V_{G+1}}$	$\frac{\partial P_N}{\partial V_{G+2}}$	\dots	$\frac{\partial P_N}{\partial V_N}$	$\frac{\partial P_N}{\partial V_2}$	$\frac{\partial P_N}{\partial V_5}$	N-G

$\frac{\partial Q_{G+1}}{\partial \delta_{G+1}}$	$\frac{\partial Q_{G+1}}{\partial \delta_{G+2}}$	\dots	$\frac{\partial Q_{G+1}}{\partial \delta_N}$	$\frac{\partial Q_{G+1}}{\partial V_{G+1}}$	$\frac{\partial Q_{G+1}}{\partial V_{G+2}}$	\dots	$\frac{\partial Q_{G+1}}{\partial V_N}$	$\frac{\partial Q_{G+1}}{\partial V_2}$	$\frac{\partial Q_{G+1}}{\partial V_5}$	N-G+1
$\frac{\partial Q_{G+2}}{\partial \delta_{G+1}}$	$\frac{\partial Q_{G+2}}{\partial \delta_{G+2}}$	\dots	$\frac{\partial Q_{G+2}}{\partial \delta_N}$	$\frac{\partial Q_{G+2}}{\partial V_{G+1}}$	$\frac{\partial Q_{G+2}}{\partial V_{G+2}}$	\dots	$\frac{\partial Q_{G+2}}{\partial V_N}$	$\frac{\partial Q_{G+2}}{\partial V_2}$	$\frac{\partial Q_{G+2}}{\partial V_5}$	
\vdots										
$\frac{\partial Q_N}{\partial \delta_{G+1}}$	$\frac{\partial Q_N}{\partial \delta_{G+2}}$	\dots	$\frac{\partial Q_N}{\partial \delta_N}$	$\frac{\partial Q_N}{\partial V_{G+1}}$	$\frac{\partial Q_N}{\partial V_{G+2}}$	\dots	$\frac{\partial Q_N}{\partial V_N}$	$\frac{\partial Q_N}{\partial V_2}$	$\frac{\partial Q_N}{\partial V_5}$	K=2(N-G)

$\frac{\partial Q_2}{\partial \delta_{G+1}}$	$\frac{\partial Q_2}{\partial \delta_{G+2}}$	\dots	$\frac{\partial Q_2}{\partial \delta_N}$	$\frac{\partial Q_2}{\partial V_{G+1}}$	$\frac{\partial Q_2}{\partial V_{G+2}}$	\dots	$\frac{\partial Q_2}{\partial V_N}$	$\frac{\partial Q_2}{\partial V_2}$	$\frac{\partial Q_2}{\partial V_5}$	2(N-G+1)
$\frac{\partial Q_5}{\partial \delta_{G+1}}$	$\frac{\partial Q_5}{\partial \delta_{G+2}}$	\dots	$\frac{\partial Q_5}{\partial \delta_N}$	$\frac{\partial Q_5}{\partial V_{G+1}}$	$\frac{\partial Q_5}{\partial V_{G+2}}$	\dots	$\frac{\partial Q_5}{\partial V_N}$	$\frac{\partial Q_5}{\partial V_2}$	$\frac{\partial Q_5}{\partial V_5}$	K'
$\frac{\partial Q_5}{\partial \delta_{G+1}}$	$\frac{\partial Q_5}{\partial \delta_{G+2}}$	\dots	$\frac{\partial Q_5}{\partial \delta_N}$	$\frac{\partial Q_5}{\partial V_{G+1}}$	$\frac{\partial Q_5}{\partial V_{G+2}}$	\dots	$\frac{\partial Q_5}{\partial V_N}$	$\frac{\partial Q_5}{\partial V_2}$	$\frac{\partial Q_5}{\partial V_5}$	K'

Compute the $K' \times j'$ matrix [B'] as follows.



for all the generator buses and thus from the vectors $[Q_{\min}]$ and $[Q_{\max}]$ compare $[Q]$ with $[Q_{\min}]$ and $[Q_{\max}]$ and check for limit violation. Temporarily include all the Q_i 's which violate the limits in the vector $[a]$, forming the augmented vector $[a']$. Include the violated limits themselves in the vector $[s]$ forming the augmented vector $[s']$. Accordingly regroup the state variables $[x]$ and $[y]$ as $[x']$ and $[y']$. For example, suppose it is found that $Q_2 < Q_2 \min$ and $Q_5 > Q_5 \max$ are two of the elements in $[Q]$ have limit violation. Then temporarily define the various vectors and indices as follows.

$$k' = k + 2 ; \quad j' = j - 2$$

$$[x'] = \begin{matrix} \delta_{G+1} \\ \delta_{G+2} \\ \vdots \\ \delta_N \\ V_{G+1} \\ V_{G+2} \\ \vdots \\ V_N \\ \dots \\ V_2 \\ V_5 \end{matrix}$$

$K' \times 1$

$$[y'] = \begin{matrix} \delta_2 \\ \delta_3 \\ \vdots \\ \delta_G \\ V_1 \\ V_3 \\ V_4 \\ V_6 \\ V_7 \\ \vdots \\ \vdots \\ \vdots \\ V_G \end{matrix}$$

$j' \times 1$

[B'] =

$\frac{\partial P_{G+1}}{\partial \delta_2}$	$\frac{\partial P_{G+1} \dots \partial P_{G+1}}{\partial \delta_3}$	$\frac{\partial P_{G+1}}{\partial \delta_G}$	$\frac{\partial P_{G+1}}{\partial v_1}$	$\frac{\partial P_{G+1}}{\partial v_3}$	$\frac{\partial P_{G+1}}{\partial v_4}$	$\frac{\partial P_{G+1}}{\partial v_6}$	$\frac{\partial P_{G+1}}{\partial v_7}$	$\frac{\partial P_{G+1}}{\partial v_G}$	1
$\frac{\partial P_{G+2}}{\partial \delta_2}$	$\frac{\partial P_{G+2} \dots \partial P_{G+2}}{\partial \delta_3}$	$\frac{\partial P_{G+2}}{\partial \delta_G}$	$\frac{\partial P_{G+2}}{\partial v_1}$	$\frac{\partial P_{G+2}}{\partial v_3}$	$\frac{\partial P_{G+2}}{\partial v_4}$	$\frac{\partial P_{G+2}}{\partial v_6}$	$\frac{\partial P_{G+2}}{\partial v_7}$	$\frac{\partial P_{G+2}}{\partial v_G}$	2
\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	
$\frac{\partial P_N}{\partial \delta_2}$	$\frac{\partial P_N \dots \partial P_N}{\partial \delta_3}$	$\frac{\partial P_N}{\partial \delta_G}$	$\frac{\partial P_N}{\partial v_1}$	$\frac{\partial P_N}{\partial v_3}$	$\frac{\partial P_N}{\partial v_4}$	$\frac{\partial P_N}{\partial v_6}$	$\frac{\partial P_N}{\partial v_7}$	$\frac{\partial P_N}{\partial v_G}$	N-G

$\frac{\partial Q_{G+1}}{\partial \delta_2}$	$\frac{\partial Q_{G+1} \dots \partial Q_{G+1}}{\partial \delta_3}$	$\frac{\partial Q_{G+1}}{\partial \delta_G}$	$\frac{\partial Q_{G+1}}{\partial v_1}$	$\frac{\partial Q_{G+1}}{\partial v_3}$	$\frac{\partial Q_{G+1}}{\partial v_4}$	$\frac{\partial Q_{G+1}}{\partial v_6}$	$\frac{\partial Q_{G+1}}{\partial v_7}$	$\frac{\partial Q_{G+1}}{\partial v_G}$	N-G+1
$\frac{\partial Q_{G+2}}{\partial \delta_2}$	$\frac{\partial Q_{G+2} \dots \partial Q_{G+2}}{\partial \delta_3}$	$\frac{\partial Q_{G+2}}{\partial \delta_G}$	$\frac{\partial Q_{G+2}}{\partial v_1}$	$\frac{\partial Q_{G+2}}{\partial v_3}$	$\frac{\partial Q_{G+2}}{\partial v_4}$	$\frac{\partial Q_{G+2}}{\partial v_6}$	$\frac{\partial Q_{G+2}}{\partial v_7}$	$\frac{\partial Q_{G+2}}{\partial v_G}$	N-G+2
\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	
$\frac{\partial Q_N}{\partial \delta_2}$	$\frac{\partial Q_N \dots \partial Q_N}{\partial \delta_3}$	$\frac{\partial Q_N}{\partial \delta_G}$	$\frac{\partial Q_N}{\partial v_1}$	$\frac{\partial Q_N}{\partial v_3}$	$\frac{\partial Q_N}{\partial v_4}$	$\frac{\partial Q_N}{\partial v_6}$	$\frac{\partial Q_N}{\partial v_7}$	$\frac{\partial Q_N}{\partial v_G}$	2(N-G)
$\frac{\partial Q_2}{\partial \delta_2}$	$\frac{\partial Q_2 \dots \partial Q_2}{\partial \delta_3}$	$\frac{\partial Q_2}{\partial \delta_G}$	$\frac{\partial Q_2}{\partial v_1}$	$\frac{\partial Q_2}{\partial v_3}$	$\frac{\partial Q_2}{\partial v_4}$	$\frac{\partial Q_2}{\partial v_6}$	$\frac{\partial Q_2}{\partial v_7}$	$\frac{\partial Q_2}{\partial v_G}$	2(N-G)H
$\frac{\partial Q_5}{\partial \delta_2}$	$\frac{\partial Q_5 \dots \partial Q_5}{\partial \delta_3}$	$\frac{\partial Q_5}{\partial \delta_G}$	$\frac{\partial Q_5}{\partial v_1}$	$\frac{\partial Q_5}{\partial v_3}$	$\frac{\partial Q_5}{\partial v_4}$	$\frac{\partial Q_5}{\partial v_6}$	$\frac{\partial Q_5}{\partial v_7}$	$\frac{\partial Q_5}{\partial v_G}$	K'
1	2	G-1	G	G+1					

Compute the $m \times K'$ matrix $[C']$ as follows:

$[c'] =$	$\frac{\partial P_1}{\partial \delta_{G+1}}$	$\frac{\partial P_1}{\partial \delta_{G+2}}$	$\frac{\partial P_1}{\partial \delta_N}$	$\frac{\partial P_1}{\partial V_{G+1}}$	$\frac{\partial P_1}{\partial V_{G+2}}$	$\frac{\partial P_1}{\partial V_N}$	$\frac{\partial P_1}{\partial V_2}$	$\frac{\partial P_1}{\partial V_5}$	1
	$\frac{\partial P_2}{\partial \delta_{G+1}}$	$\frac{\partial P_2}{\partial \delta_{G+2}}$	$\frac{\partial P_2}{\partial \delta_N}$	$\frac{\partial P_2}{\partial V_{G+1}}$	$\frac{\partial P_2}{\partial V_{G+2}}$	$\frac{\partial P_2}{\partial V_N}$	$\frac{\partial P_2}{\partial V_2}$	$\frac{\partial P_2}{\partial V_5}$	2
	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮	⋮
	$\frac{\partial P_G}{\partial \delta_{G+1}}$	$\frac{\partial P_G}{\partial \delta_{G+2}}$	$\frac{\partial P_G}{\partial \delta_N}$	$\frac{\partial P_G}{\partial V_{G+1}}$	$\frac{\partial P_G}{\partial V_{G+2}}$	$\frac{\partial P_G}{\partial V_N}$	$\frac{\partial P_G}{\partial V_2}$	$\frac{\partial P_G}{\partial V_5}$	G
	0	0	0	0	0	0	0	0	G+1
	0	0	0	0	0	0	1	0	G+2
	0	0	0	0	0	0	0	0	⋮
	0	0	0	0	0	0	0	0	⋮
	0	0	0	0	0	0	0	0	⋮
	0	0	0	0	0	0	0	0	2G
0	0	0	1	0	0	0	0	2G+1	
0	0	0	0	1	0	0	0	⋮	
0	0	0	0	0	1	0	0	m=N+G	
1	2	N-G	N-G+1	N-G+2	K=	K+1	K	K'	

$2(N-G)$

Step 5

compute $[A']^{-1}$ and store it in the same location as $[A']$ and

compute $[C'] [A']^{-1}$ and store it in the same location as $[C']$

Compute $[G'] = [D'] - [C'] [A']^{-1} [B]$ and store it in the same location as $[D']$

Compute $[A']^{-1} [B']$ and store it in the same location as $[B']$

Compute $[H] = [G']^{\text{tr}} [W]$ where $[H]$ is a $J' \times m$ matrix

Compute $[H] [G']$ and store it in the upper part of the same location as $[G']$

Compute $([H] [G'])^{-1}$ and store it in the same location as $[H] [G']$

Compute $[A']^{-1} ([S'] - [a'])$ and store it in the same location as $[S'] - [a']$

Compute $[t] = [b] - [c'] [A']^{-1} ([s'] - [a])$ and store it in the same location as $[t] - [b]$

Compute $[\Delta y'] = ([H] [G'])^{-1} [H] ([t] - [b] - [c'] [A']^{-1} ([s'] - [a']))$

from equation 7

Compute $[\Delta x'] = [A']^{-1} ([s'] - [a'] - [A']^{-1} [B'] [\Delta y'])$

From equation 6

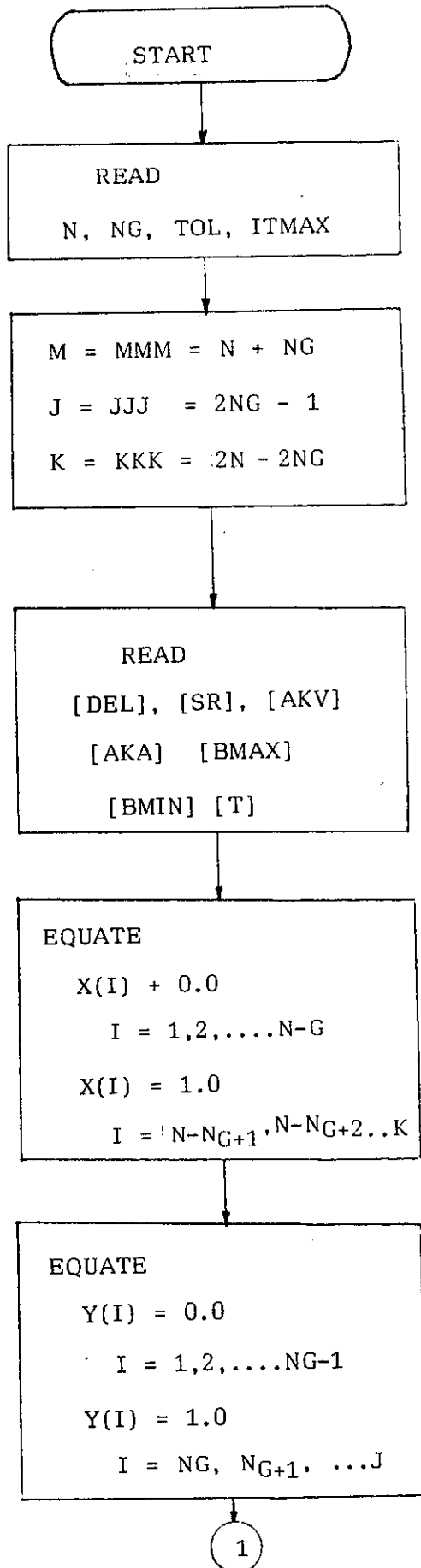
Step 6

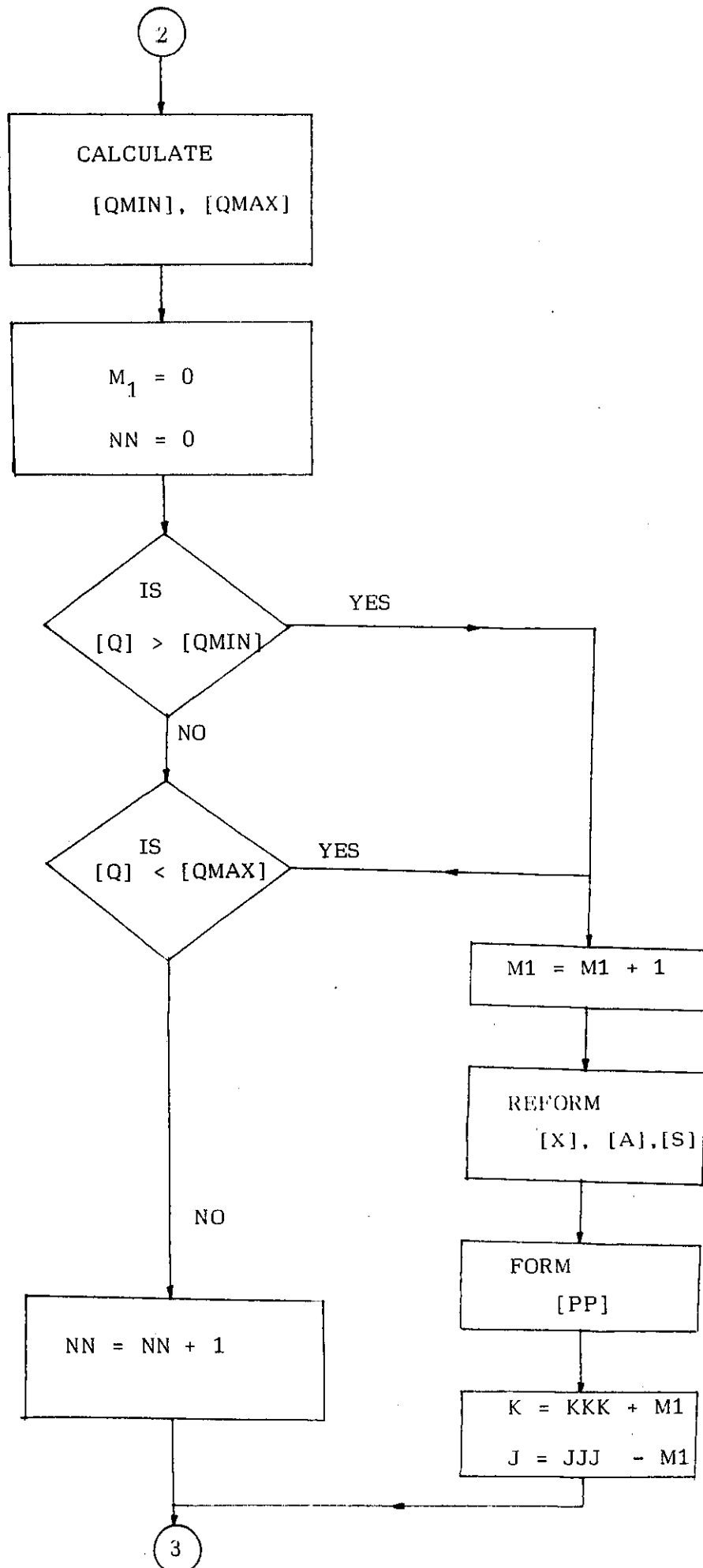
Check whether all the elements in $[\Delta x']$ and $[\Delta y']$ are less than a pre specified ϵ in magnitude of so, the solution has converged. If not go to step 7.

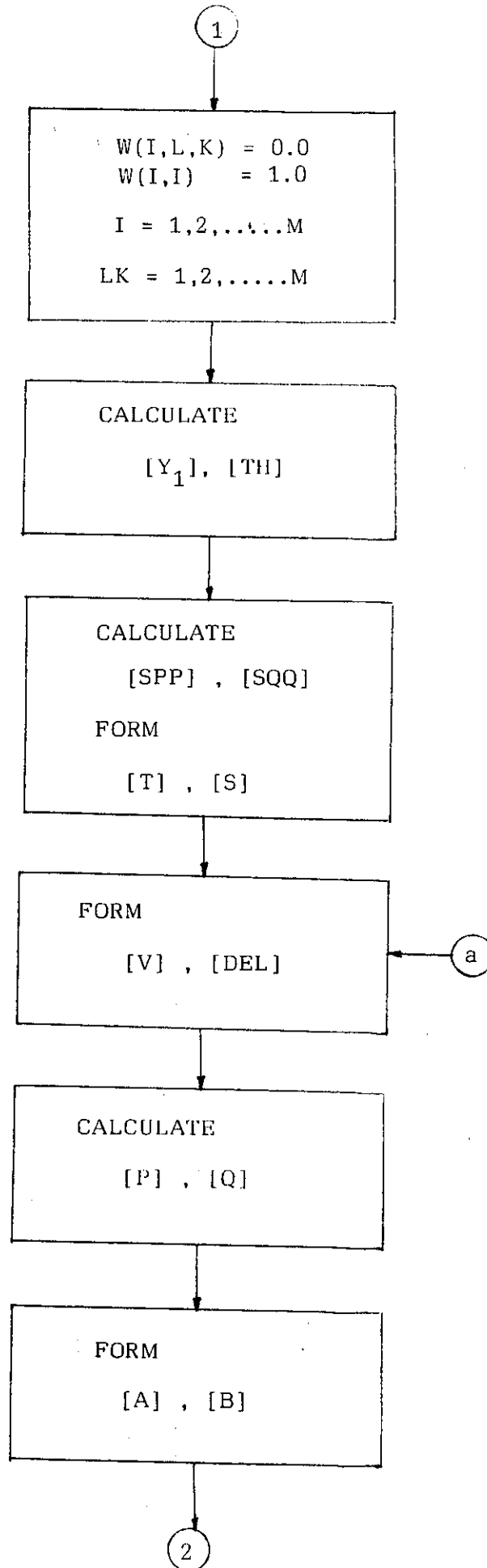
flow diagram

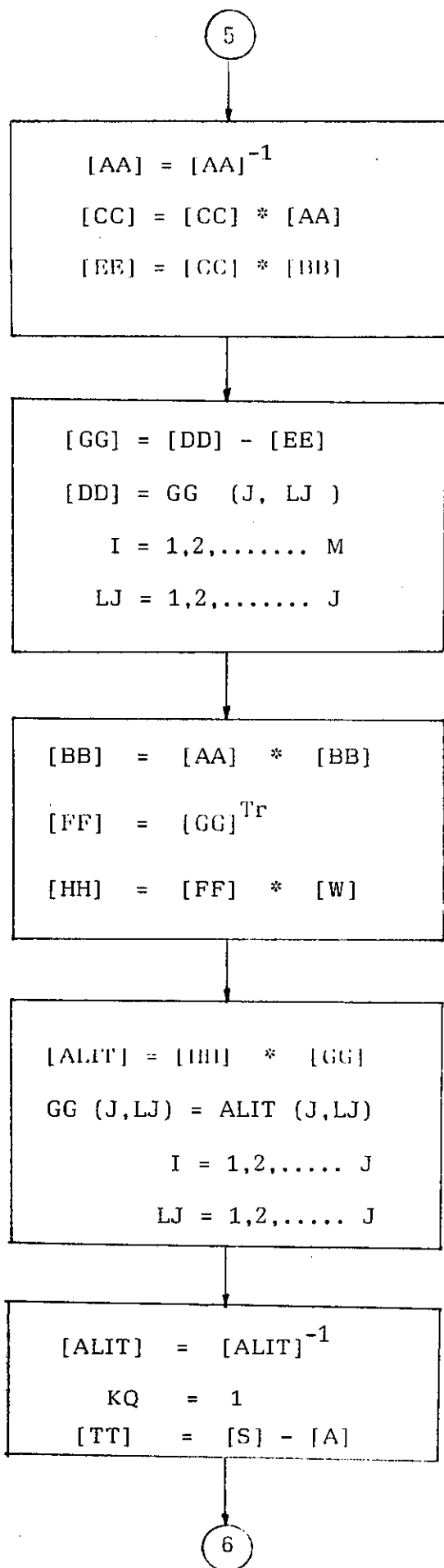
CHAPTER VI

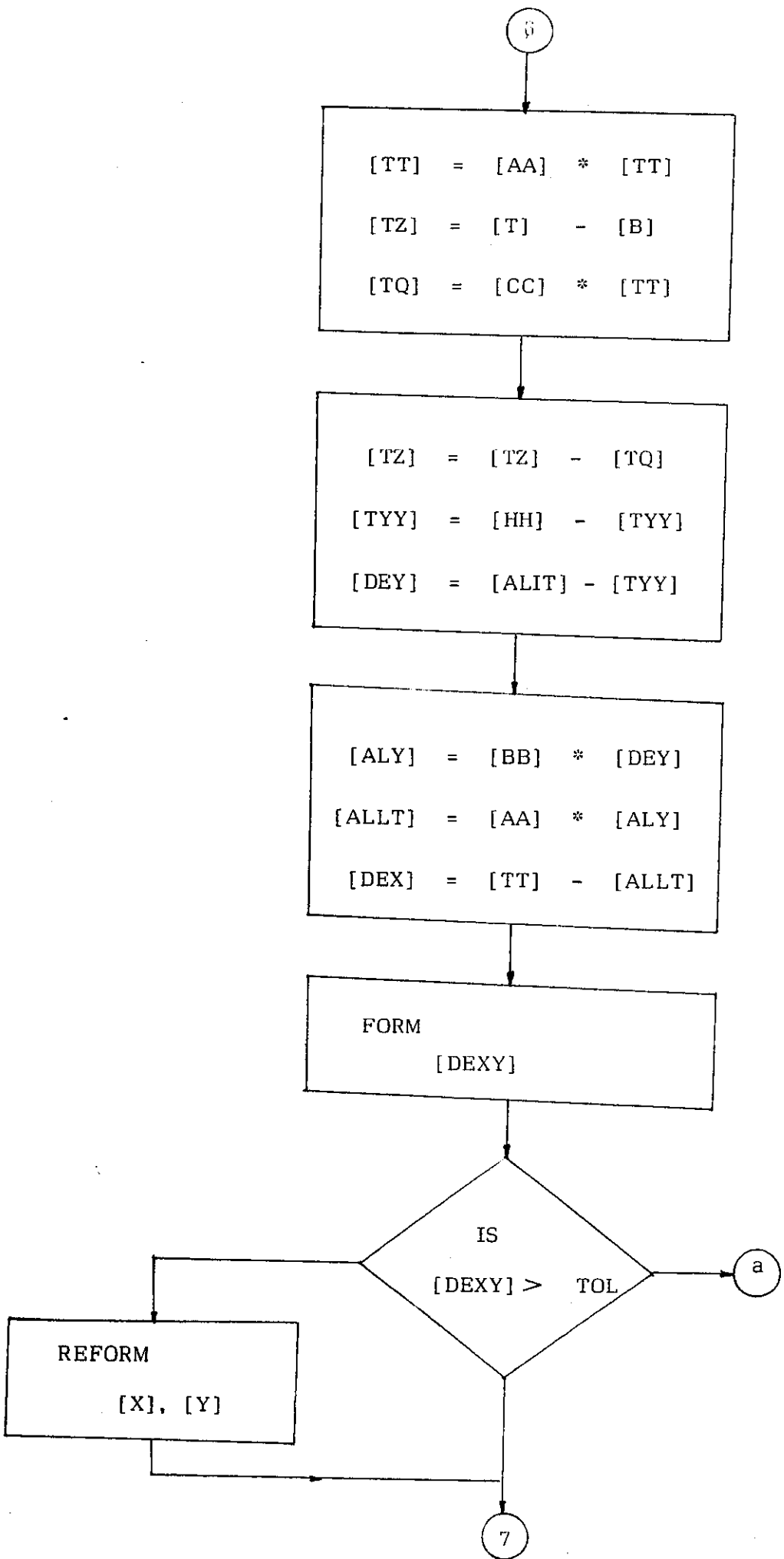
FLOW DIAGRAM:

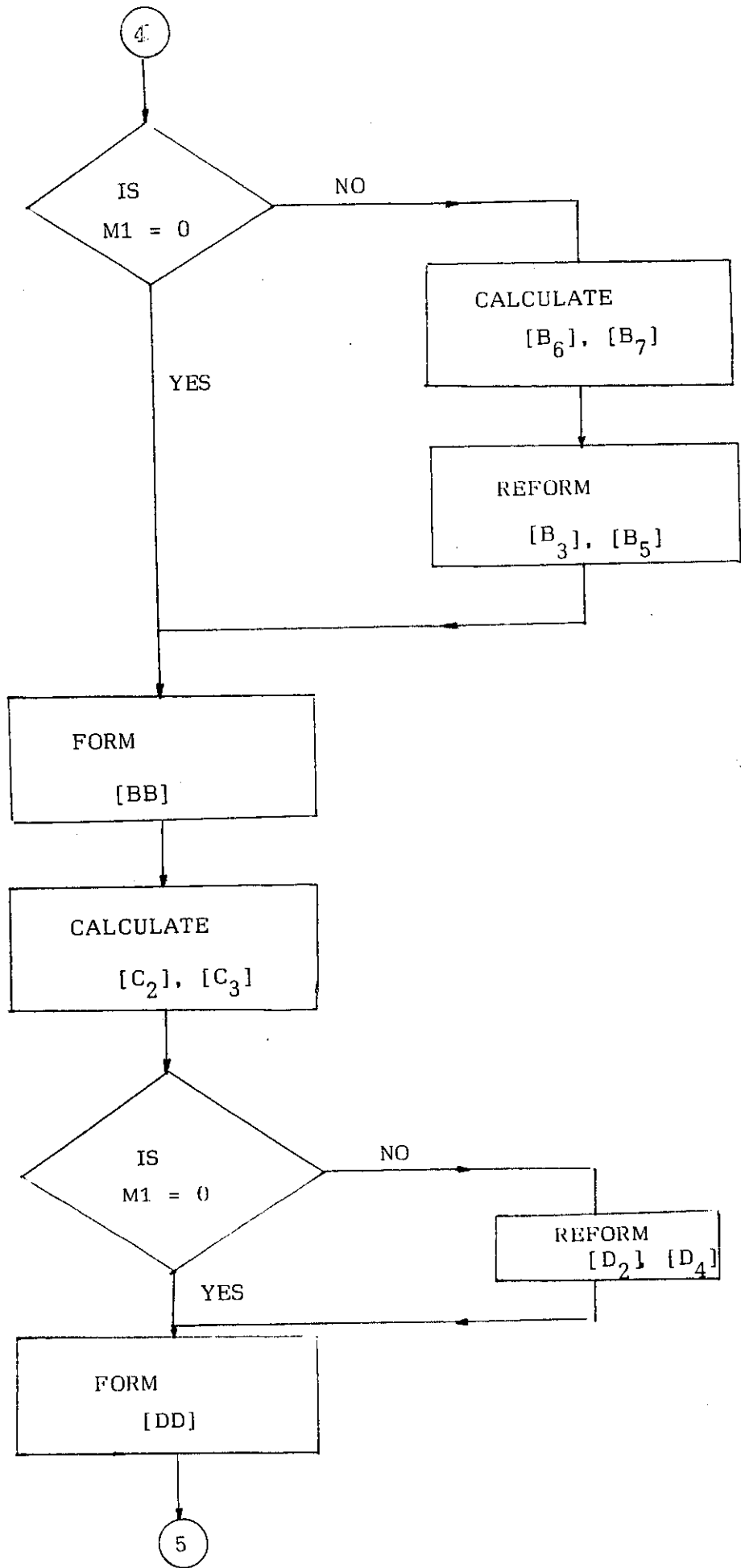


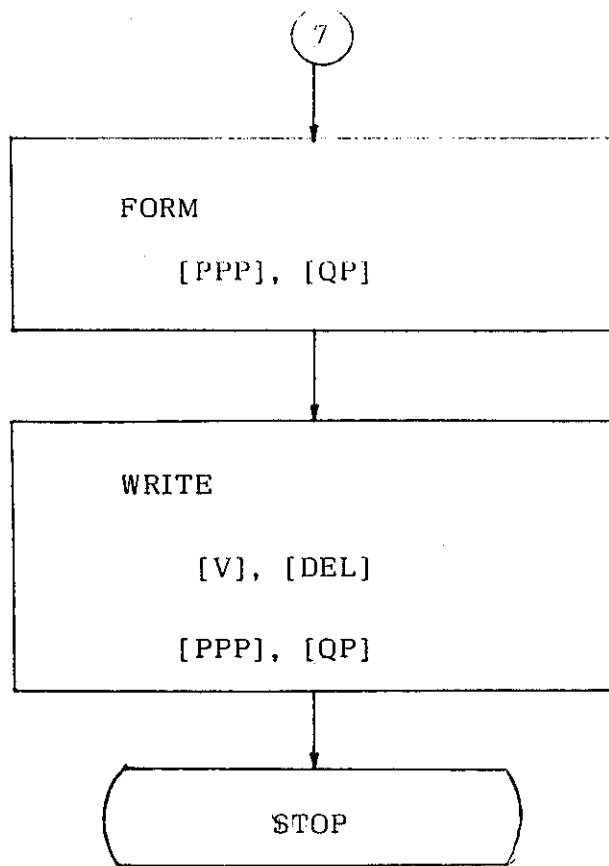












software and results


```

DIMENSION DEL(5),SR(2),ASV(5),AKA(3),BMAX(7),GMIN(7),CZ(10,10)
DIMENSION C2(10,10),S1(8,8),BP(5),PPP(5),ALLT(8,8),A6(3,2)
DIMENSION CH(10,10),I(8,8),X(8),Y(5),U(8,8),SPP(5),S(8,8)
DIMENSION A2(3,3),A8(2,3),A7(3,2),A9(2,3),A10(2,2),TZ(8,8)
DIMENSION AA(8,8),B2(2,1),B4(3,1),B3(3,2),B5(3,2),B6(2,1)
DIMENSION B7(2,2),B8(8,8),L2(2,3),C5(2,3),C6(2,2),C7(2,2)
DIMENSION C9(3,3),CC(8,8),D1(2,1),D2(2,2),D4(2,2),DEXY(8,8)
DIMENSION EE(8,8),GG(8,8),FF(3,3),HH(8,8),ALIT(8,8),TT(8,8)
DIMENSION TR(8,3),TYY(8,8),DEY(8,8),DEX(8,8),Y1(10,10)
DIMENSION ALY(8,8),SOO(5),PP(2),A3(3,3),A4(3,3),A5(3,3)
DIMENSION IDY(3),V(3),A(8,8),P(5),B(8,8),QMIN(2),QMAX(2),Q(5)
DIMENSION B1(8,8),TH(10,10),BD(8,8),CA(10,10),CY(10,10)
COMPLEX CH,CZ,CY,CA
OPEN(2,FILE='S.DAT',STATUS='OLD')

```

C
C
C

```

#####
#####          MAIN          PROGRAM          #####
#####

```

```

READ(2,199) N,NG,IUI,IIBS
M=N+NG
J=2*NG-1
K=2*(N-NG)
MM=M
JJ=J

```

KKK=K

READ(2,200) DEL(1)

READ(2,201) (SR(I), I=1,N)

READ(2,202) (SKV(I), I=1,N)

READ(2,202) (AKA(I), I=1,N)

READ(2,203) (BNAX(I), I=1,N)

READ(2,203) (BMIN(I), I=1,N)

DO 1 I=1,N

READ(2,204) (CZ(I,LJ), LJ=1,N)

1 CONTINUE

DO 2 I=1,N

READ(2,204) (CH(I,LJ), LJ=1,N)

2 CONTINUE

READ(2,205) (T(I,1), I=NG+1,N)

C #####
C [X] : DEPENDENT VARIABLES #####
C #####

DO 3 I=1, KKK

IF(I.GE.(N-NU+1)) GOTO 4

X(I)=0.0

GOTO 3

4 X(I)=1.0

3 CONTINUE

C #####
C [Y] : INDEPENDENT VARIABLES #####
C #####

DO 5 I=1, JJJ

IF(I.GE.NG) GOTO 6

Y(I)=0.0

GOTO 5

6 Y(I)=1.0

5 CONTINUE

[W] : WEIGHTAGES #####
#####

DO 7 I=1,M

DO 8 LK=1,N

IF(I.EQ.LK) GOTO 9

W(I,LK)=0.0

GOTO 8

9 W(I,LK)=1.0

8 CONTINUE

7 CONTINUE

CALL YBUS(CZ,CH,Y1,TH,N)

CALL NR (AKV,AKA,Y1,TH,N,SPP,SMO)

[T] : TARGET VALUE #####
#####

DO 10 I=1,NG

T(I,1)=SPP(I)

10 CONTINUE

[S] : SPECIFIED VALUES #####
#####

DO 11 I=1,KKK

IF(I.GE.(N-NG+1)) GOTO 12

S(I,1)=SPP(NG+I)

GOTO 11

C
C
C

```
#####
## LOAD FLOW AS A PLANNING TOOL - AN OPTIMAL APPROACH ##
#####
```

```
WRITE(*,*) ' LOAD FLOW AS A PLANNING TOOL
WRITE(*,*) ' - AN OPTIMAL APPROACH
```

WRITE(*,*)

WRITE(*,*)

WRITE(*,*) ' WORK DONE BY:

WRITE(*,*) ' R.SHANTHI

WRITE(*,*) ' M.JAYA KUMAR

WRITE(*,*) ' V.S.RAVICHANDRAN

WRITE(*,*) ' G.SHANTHANAYAKI

WRITE(*,*) ' SHINY PURUSHOTHAMAN

WRITE(*,*)

WRITE(*,*)

WRITE(*,*) ' GUIDED BY:

WRITE(*,*) ' MR.K.SURENDRAN M.Tech.,

S(I,1)=SUG(I-N+2*NG)

11 CONTINUE

ITERATION STARTS #####
#####

DO 2000 IAB=1,ITMAX

DO 13 I=1,JJJ

13 IDY(I)=1

[V] : VOLTAGE MAGNITUDE #####
[DEL]: VOLTAGE ANGLE IN RADIANS #####
#####

DO 14 I=1,N

IF(I.GT.NG) GOTO 15

V(I)=Y(I+NG-1)

IF(I.EQ.1) GOTO 16

DEL(I)=Y(I-1)

16 GOTO 14

15 V(I)=X(I+N-2*NG)

DEL(I)=X(I-NG)

14 CONTINUE

CALL NR(V,DEL,Y1,TH,N,P,Q)

[A] : SPECIFIED VARIABLES #####
#####

DO 17 I=1,N-NG

17 A(I,1)=P(NG+I)

DO 18 I=N-NG+1,KKK

18 A(I,1)=Q(2*NG+I-N)

C

C
C

[B] : TARGET VARIABLES #####
#####

DO 19 I=1,NG

19 B(I,1)=F(I)

DO 20 I=NG+1,MMN

20 B(I,1)=V(I-NG)

C
C
C
C

[QMAX] : MAXIMUM VALUE OF REACTIVE POWER #####
[QMIN] : MINIMUM VALUE OF REACTIVE POWER #####
#####

DO 21 I=1,NG

QMIN(I)=-SQRT(SR(I)*SR(I)-P(I)*P(I))

21 QMAX(I)=ABS(QMIN(I))

C
C
C

CHECK FOR REACTIVE POWER VIOLATION #####
#####

M1=0

NN=0

DO 22 I=1,NG

IF(Q(I).GT.QMIN(I)) GOTO 23

M1=M1+1

X(KKK+M1)=V(I)

A(KKK+M1,1)=Q(I)

IDY(I+NG-1)=0

PP(I-NN)=I

S(KKK+M1,1)=QMIN(I)

K=KKK+M1

J=JJJ-M1

GOTO 22

23 IF(Q(I).LT.QMAX(I)) GOTO 24

M1=M1+1

```

33      CONTINUE
        DO 35 I=1,M1
        DO 36 KJ=1,M1

        JI1=PP(I)
        JI2=PP(KJ)

        IF(JI1.EQ.JI2) GOTO 37

        THY1 = TH(JI1,JI2)+DEL(JI1)-DEL(JI2)

        A10(I,KJ)=ABS(V(JI1)*Y1(JI1,JI2))*SIN(THY1)

        GOTO 36
37      A10(I,KJ)=0

        DO 38 LK=1,N

        IF(LK.EQ.JI1) GOTO 38

        THYI = TH(JI1,LK)+DEL(JI1)-DEL(LK)

        A10(I,KJ)=A10(I,KJ)+ABS(V(LK)*Y1(JI1,LK))*SIN( THYI )

38      CONTINUE

        A10(I,KJ)=A10(I,KJ)+2*ABS(V(JI1)*Y1(JI1,JI1))*SIN(TH(JI1,JI1))

36      CONTINUE

35      CONTINUE

C      #####
C      ##### FORMATION OF [AA] MATRIX #####
C      #####

32      DO 40 I=1,N-NG

        DO 40 LJ=1,N-NG

40      AA(I,LJ)=A2(I,LJ)

        DO 41 I=1,N-NG

        DO 41 LJ=N-NG+1,KKK

41      AA(I,LJ)=A3(I,LJ-N+NG)

        DO 42 I=N-NG+1,KKK

        DO 42 LJ=1,N-NG

```

```

A5(IT,IL)=ABS(V(I)*Y1(I,LJ))*SIN(TH(I,LJ)+DEL(I)-DEL(LJ))
GOTO 29
30  A2(IT,IL)=0.0
    A3(IT,IL)=0.0
    A4(IT,IT)=0.0
    A5(IT,IL)=0.0
    DO 31 LL=1,N
    IF(LL.EQ.1) GOTO 31
    THT = TH(I,LL)+DEL(I)-DEL(LL)
    A2(IT,IL)=-ABS(V(I)*V(LL)*Y1(I,LL))*SIN(THT)+A2(IT,IL)
    A3(IT,IL)=ABS(V(LL)*Y1(I,LL)) * COS(THT) + A3(IT,IL)
    A4(IT,IL)=ABS(V(I)*V(LL)*Y1(I,LL))*COS(THT) + A4(IT,IL)
    A5(IT,IL)=ABS(V(LL)*Y1(I,LL))*SIN(THT)+A5(IT,IL)
31  CONTINUE
    A3(IT,IL)=2*ABS(V(I)*Y1(I,1))*COS(TH(I,1))+A3(IT,IL)
    A5(IT,IL)=2*ABS(V(I)*Y1(I,1))*SIN(TH(I,1))+A5(IT,IL)
29  CONTINUE
28  CONTINUE
    IF(M1.EQ.0) GOTO 32
    DO 33 I=1,M1
    NN1=PP(I)
    DO 34 K1=NG+1,N
    THH = TH(K1,NN1)+DEL(K1)-DEL(NN1)
    THY = TH(NN1,K1)+DEL(NN1)-DEL(K1)
    A6(K1-NG,I)=ABS(V(K1)*Y1(K1,NN1))*COS(THH)
    A8(I,K1-NG)=-ABS(V(K1)*V(NN1)*Y1(NN1,K1))*COS(THY)
    A7(K1-NG,I)=ABS(V(K1)*Y1(K1,NN1))*SIN(THH)
    A9(I,K1-NG)=ABS(V(NN1)*Y1(NN1,K1))*SIN(THY)
34  CONTINUE

```



```

X(KKK+M1)=V(I)
A(KKK+M1,1)=Q(I)
IDY(NG-1+I)=0
PF(I-NN)=I
S(KKK+M1,1)=QMAX(I)
K=KKK+M1
J=JJJ-M1

```

```
GOTO 22
```

```
24 NN=NN+1
```

```
22 CONTINUE
```

```

C #####
C ##### CHECK FOR TARGET VIOLATION #####
C #####

```

```
DO 25 I=1,MMI4
```

```
IF(B(I,1).GT.BMIN(I).AND.B(I,1).LT.BMAX(I)) GOTO 25
```

```
W(I,I)=W(I,I)+1.0
```

```
25 CONTINUE
```

```

C #####
C ##### FORMATION OF [A2],[A3],[A4],[A5],[A6],[A7],[A8],[A9],[A10]#
C #####

```

```
DO 28 I=NG+1,N
```

```
DO 29 LJ=NG+1,N
```

```
IT=I-NG
```

```
IL=LJ-NG
```

```
IF(I.EQ.LJ) GOTO 30
```

```
THIP=TH(I,LJ)+DEL(I)-DEL(LJ)
```

```
A2(IT,IL)=ABS(V(I)*V(LJ)*Y1(I,LJ))*SIN(THIP)
```

```
A3(IT,IL)=ABS(V(I)*Y1(I,LJ))*COS(THIP)
```

```
THI1=TH(I,LJ)+DEL(I)-DEL(LJ)
```

```
A4(IT,IL)=-ABS(V(I)*V(LJ)*Y1(I,LJ))*COS(THI1)
```

```

42      AA(I,LJ)=A4(I-N+NG,LJ)
      DO 43 I=N-NG+1,KKK
      DO 43 LJ=N-NG+1,KKK
43      AA(I,LJ)=A5(I-N+NG,LJ-N+NG)
      IF(M1.EQ.0) GO TO 44
      DO 45 I=1,N-NG
      DO 45 LJ=KKK+1,K

```

```

45      AA(I,LJ)=A6(I,LJ-KKK)
      DO 46 I=N-NG+1,KKK
      DO 46 LJ=KKK+1,K

```

```

46      AA(I,LJ)=A7(I-N+NG,LJ-KKK)
      DO 47 I=KKK+1,K
      DO 47 LJ=1,N-NG

```

```

47      AA(I,LJ)=A8(I-KKK,LJ)
      DO 48 I=KKK+1,K
      DO 48 LJ=N-NG+1,KKK

```

```

48      AA(I,LJ)=A9(I-KKK,LJ-N+NG)
      DO 49 I=KKK+1,K
      DO 49 LJ=KKK+1,K

```

```

49      AA(I,LJ)=A10(I-KKK,LJ-K)

```

```

C      #####
C      #####          FORMATIN OF [B2],[B3],[B4],[B5],[B6],[B7] #####
C      #####

```

```

44      DO 50 I=NG+1,N
      DO 51 LJ=1,NG
      LG1=I-NG
      UHI=TH(I,LJ)+DEL(I)-DEL(LJ)
      IF(LJ.EQ.1) GOTO 52
      B2(LG1,LJ-1)=ABS(V(I)*V(LJ)*Y1(I,LJ))*SIN(UHI)
      B4(LG1,LJ-1)=-ABS(V(I)*V(LJ)*Y1(I,LJ))*COS(UHI)

```

```

52     B3(LG1,LJ)=ABS(V(I)*Y1(I,LJ))*COS(UH1)
      B5(LG1,LJ)=ABS(V(I)*Y1(I,LJ))*SIN(UH1)
51     CONTINUE
50     CONTINUE
      IF(M1.EQ.0) GOTO 53
      DO 54 I=1,M1
      NV=PP(I)
      DO 54 IJK=1,N-NG
      B3(IJK,NV)=B3(IJK,NV+1)
54     B5(IJK,NV)=B5(IJK,NV+1)
      DO 55 II=1,M1
      NA=PP(II)
      MJ=0
      DO 56 LJ=1,NG
      IF(NA.EQ.LJ) GOTO 57
      HIJ=TH(NA,LJ)+DEL(NA)-DEL(LJ)
      B7(I,LJ-MJ)=ABS(V(NA)*Y1(NA,LJ))*SIN(HIJ)
      GOTO 56
57     MJ=MJ+1
56     CONTINUE
55     CONTINUE
      DO 58 IK=1,M1
      NA=PP(I)
      DO 59 JK=2,NG
      IF(NA.EQ.J) GOTO 60
      THY2 = TH(NA,LJ)+DEL(NA)-DEL(LJ)
      B6(I,LJ-1)=-ABS(V(NA)*Y1(NA,LJ))*COS(THY2)
      GOTO 59

```

```

60   B6(NA,LJ-1)=0
      DO 610 KK=1,N
      IF(NA.EQ.KK) GOTO 610
      THUY = TH(NA, KK)+DEL(NA)-DEL(KK)
      B6(NA,LJ-1)=ABS(V(NA)*V(KK)*Y1(NA, KK))*COS(THUY) + B6(NA,LJ-1)

```

```

610  CONTINUE
59   CONTINUE
58   CONTINUE

```

```

C   #####
C   ##### FORMATION OF [BB] MATRIX #####
C   #####

```

```

53   DO 600 I = 1,N-NG
      DO 606 JJ = 1,NG-1
      BB(I,JJ)=B2(I,JJ)
606  CONTINUE
600  CONTINUE
      DO 61 I=1,N-NG
      DO 61 JJ=NG,J
61   BB(I,JJ)=B3(I,JJ-NG+1)
      DO 62 I=N-NG+1,KKK
      DO 62 JJ=1,NG-1
62   BB(I,JJ)=B4(I-N+NG,JJ)
      DO 63 I=N-NG+1,KKK
      DO 63 JJ=NG,J
63   BB(I,JJ)=B5(I-N+NG,JJ-NG+1)
      IF(M1.EQ.0) GOTO 64
      DO 65 I=KKK,K
      DO 65 JJ=1,NG-1
65   BB(I,JJ)=B6(I-K+1,JJ)
      DO 66 I=KKK,K

```

```

DO 82 LJ=N-NG+1, KKK
IG=I-2*NG
IJ1=LJ-N+NG
82 CC(I, LJ)=C9(IG, IJ1)
IF(M1.EQ.0) GOTO 83
DO 84 I=1, NG
DO 84 LJ=KKK+1, K
84 CC(I, LJ)=D6(I, LJ-KKK)
DO 85 I=NG+1, 2*NG
DO 85 LJ=KKK+1, K
85 CC(I, LJ)=C7(I-NG, LJ-KKK)
DO 86 I=2*NG+1, MMM
DO 86 LJ=KKK+1, K
86 CC(I, LJ)=0.0
C #####
C ##### FORMATION OF [D1], [D2], [D3], [D4], [D5], [D6] #####
C #####
83 DO 87 I=1, NG
DO 88 LJ=1, NG
IF(LJ.EQ.1) GOTO 89
D4(I, LJ)=0.0
D2(I, LJ)=ABS(V(I)*Y1(I, LJ))*COS(TH(I, LJ)+DEL(I)-DEL(LJ))
IF(LJ.EQ.1) GOTO 91
OLF=TH(I, LJ)+DEL(I)-DEL(LJ)
D1(I, LJ-1)=ABS(V(I)*V(LJ))*Y1(I, LJ)*SIN(OLF)
91 GOTO 88
89 D1(I, LJ-1)=0.0
D2(I, LJ)=0.0
D4(I, LJ)=1.0
DO 90 LK=1, N

```

```

GOTO 71
72  C7(KJ,NK)=1.0
    C6(KJ,NK)=0.0
    DO 73 LK=1,N
    IF(KJ.EQ.LK) GOTO 73
    TKLI = TH(KJ,LK)+DEL(KJ)-DEL(LK)
    C6(KJ,NK)=ABS(V(LK)*Y1(KJ,LK))*COS(TKLI)+C6(KJ,NK)
73  CONTINUE
    C6(KJ,NK)=2*ABS(V(KJ)*Y1(KJ,KJ))*COS(TH(KJ,KJ))+C6(KJ,NK)
71  CONTINUE
70  CONTINUE
C   #####
C   ##### FORMATION OF [CC] #####
C   #####
69  DO 77 I=1,NG
    DO 77 LJ=1,N-NG
77  CC(I,LJ)=C2(I,LJ)
    DO 78 I=1,NG
    DO 78 LJ=N-NG+1,KKK
78  CC(I,LJ)=C3(I,LJ-N+NG)
    DO 79 I=NG+1,2*NG
    DO 79 LJ=1,N-NG
79  CC(I,LJ)=0.0
    DO 80 I=NG+1,2*NG
    DO 80 LJ=N-NG+1,KKK
80  CC(I,LJ)=0.0
    DO 81 I=2*NG+1,MMM
    DO 81 LJ=1,N-NG
81  CC(I,LJ)=0.0
    DO 82 I=2*NG+1,MMM

```

IF(I.EQ.LK) GOTO 90

TH2 =TH(I,LK)+DEL(I)-DEL(LK)

D2(I,LJ)=ABS(V(LK)*Y1(I,LK))*COS(TH2)+D2(I,LJ)

IF(LJ.EQ.1) GOTO 90

THUI= TH(I,LK)+DEL(I)-DEL(LK)

D1(I,LJ-1)=-ABS(V(I)*V(LK)*Y1(I,LK))*SIN(THUI)+D1(I,LJ-1)

90 CONTINUE

D2(I,LJ)=2*ABS(V(I)*Y1(I,I))*COS(TH(I,I))+D2(I,LJ)

88 CONTINUE

87 CONTINUE

IF(M1.EQ.0) GOTO 92

DO 93 I=1,M1

NL=PP(I)

DO 94 LK = 1,NG

D2(LK,NL)=D2(LK,NL+1)

D4(LK,NL)=D4(LK,NL+1)

94 CONTINUE

93 CONTINUE

C #####
C ##### FORMATION OF [OD] #####
C #####

92 DO 95 I=1,NG

DO 95 LJ=1,NG-1

95 DD(I,LJ)=D1(I,LJ)

DO 96 I=1,NG

DO 96 LJ=NG,J

96 DD(I,LJ)=D2(I,LJ-NG+1)

DO 97 I=NG+1,2*NG

DO 97 LJ=1,NG-1

```

DO 66 JJ=NG,J
66  BB(I,JJ)=B7(I-K+1,JJ-NG+1)

#####
FORMATION OF [C2],[C3],[C4],[C5],[C6],[C7],[C8],[C9],[C10]
#####

64  DO 67 I=1,NG

DO 68 JJ=NG+1,N

J1J=JJ-NG

HIJL=TH(I,JJ)+DEL(I)-DEL(JJ)

C2(I,J1J)=ABS(V(I)*V(JJ)*Y1(I,JJ))*SIN(HIJL)

C3(I,J1J)=ABS(V(I)*Y1(I,JJ))*COS(HIJL)

68  CONTINUE

67  CONTINUE

DO 74 I=1,N-NG

DO 75 LJ=1,N-NG

IF (I.EQ.LJ) GOTO 76

C9(I,LJ)=0.0

GOTO 75

76  C9(I,LJ)=1.0

75  CONTINUE

74  CONTINUE

IF(M1.EQ.0) GOTO 69

DO 70 I=1,M1

NK=PP(I)

DO 71 KJ=1,NG

IF(KJ.EQ.NK) GOTO 72

TH1 =TH(KJ,NK)+DEL(K)-DEL(NK)

C6(KJ,NK)=ABS(V(KJ)*Y1(KJ,NK))* COS(TH1)

C7(KJ,NK)=0.0

```



```

97 DD(I,LJ)=0.0
DO 98 I=NG+1,2*NG
DO 98 LJ=NG,J
98 DD(I,LJ)=D4(I-NG,LJ-NG+1)
DO 99 I=2*NG+1,MMM
DO 99 LJ=1,J
99 DD(I,LJ)=0.0
C ##### [AA] = [AA] INVERSE #####
CALL MINE(AA,AA,K)
C ##### [CC] = [CC] X [AA] #####
CALL MTML(CC,AA,CC,M,K,K)
C ##### [EE] = [CC] X [BB] #####
CALL MTML(CC,BB,EE,M,K,J)
C ##### [GG] = [DD] - [EE] #####
CALL SUB(DD,EE,GG,M,J)
C ##### [DD] = [GG] #####
DO 100 I = 1,M
DO 100 LJ = 1,J
100 DD(I,LJ) = GG(I,LJ)
C ##### [BB] = [AA] X [BB] #####
CALL MTML(AA,BB,BB,K,K,J)
C ##### [FF] = [GG] TRANSPOSE #####
CALL TRANS(GG,FF,M,J)
C ##### [HH] = [FF] X [W] #####
CALL MTML(FF,W,HH,J,M,M)
C ##### [ALIT] = [HH] X [GG] #####
CALL MTML(HH,GG,ALIT,J,M,J)
C ##### UPPER PART [GG] CONTAINS [ALIT] #####
DO 101 I = 1,J

```

```

DO 101 LJ = 1,J
101  GG(I,LJ) = ALIT(I,LJ)
C    ##### [ALIT] = [ALIT] INVERSE #####
CALL MINE(ALIT,ALIT,J)
KQ=1
C    ##### [TT] = [S] - [A] #####
CALL SUB(S,A,TT,K,KQ)
C    ##### [TT] = [AA] X [TT] #####
CALL MTML(AA,TT,TT,K,K,KQ)
C    ##### [TZ] = [T] - [B] #####
CALL SUB(T,B,TZ,M,KQ)
C    ##### [TQ] = [CC] X [TT] #####
CALL MTML(CC,TT,TQ,M,K,KQ)
C    ##### [TZ] = [TZ] - [TQ] #####
CALL SUB(TZ,TQ,TZ,M,KQ)
C    ##### [TYT] = [HH] X [TYT] #####
CALL MTML(HH,TZ,TYT,J,M,KQ)
C    ##### [DEY] = [ALIT] X [TYT] #####
CALL MTML(ALIT,TYT,DEY,J,J,KQ)
C    ##### [ALY] = [BB] X [DEY] #####
CALL MTML(BB,DEY,ALY,K,J,KQ)
C    ##### [ALLT] = [AA] X [ALY] #####
CALL MTML(AA,ALY,ALLT,K,K,KQ)
C    ##### [DEX] = [TT] - [ALLT] #####
CALL SUB(TT,ALLT,DEX,K,KQ)
DO 102 I=1,K+J
IF(I.EQ.K) GOTO 113
DEXY(I,1)=DEX(I,1)
GOTO 103

```

```

113   DEXY(I,1)=DEY(I-K,1)
103   IF(DEXY(I,1).GT.TOL) GOTO 105
102   CONTINUE
      GOTO 106
105   DO 107 I=1,K
      X(I)=X(I)+DEX(I,1)
      IF(I.GT.J)GOTO 107
      Y(I)=Y(I)+DEY(I,1)
107   CONTINUE
      DO 108 I=1,M1
      NJ=PP(I)
      Y(NG-1+NJ)=X(KKK+I)
108   CONTINUE
2000  CONTINUE
106   DO 109 I=1,NG
      PFF(I)=B(I,1)
      QF(I)=Q(I)
109   CONTINUE
      DO 110 I=1,KKK
      IF (I.GE.N-NG+1) GOTO 111
      PFF(NG+I)=A(I,1)
      GOTO 110
111   QF(I-1)=A(I,1)
110   CONTINUE
      WRITE(*,400)
400   FORMAT(5X,'VOLTAGE',7X,'VOLTAGE ANGLE',7X,'      REAL',5X,
1     '      REACTIVE')
      WRITE(*,401)
401   FORMAT(22X,10HIN RADIANS,13X,5HPOWER,15X,5HPOWER)

```

```

DO 112 I=1,N
WRITE(*,206) V(I),DEL(I),PPP(I),BP(I)
112 CONTINUE
206 FORMAT(3X,F8.5,10X,F8.5,14X,F10.5,13X,F8.5)
199 FORMAT(I2,I2,F10.6,I2)
200 FORMAT(F8.5)
201 FORMAT(2F6.2)
202 FORMAT(5F8.5)
203 FORMAT(7F10.4)
204 FORMAT(10F7.3)
205 FORMAT(5F7.3)
500 FORMAT(/2X,8F10.5/)
STOP
END

```

```

C ##### FORMING THE YBUS #####
SUBROUTINE YBUS(CZ,CH,Y1,TH,N)
COMPLEX CH,CZ,CY,CA
DIMENSION CZ(10,10), CA(10,10), CH(10,10), CY(10,10),Y1(10,10)
DIMENSION TH(10,10)
DO 7 I=1,N
DO 7 LJ=1,N
IF (CZ(I,LJ).EQ.0.0) GO TO 7
CA(I,LJ)=1.0/CZ(I,LJ)
7 CONTINUE
DO 25 I=1,N
CY(I,I)=0.0
DO 35 LJ=1,N
IF (I.EQ.LJ) GOTO 35

```

```

LR = I
LC = J
120 CONTINUE
130 CONTINUE
    IF ( AL . EQ . 0.0 ) GO TO 270
    IF ( LR . EQ . IC . AND . LC . EQ . IC ) GOTO 160
    DO 140 I = 1,N
        D(I) = B(IC,I)
        B(IC,I) = B(LR,I)
        B(LR,I) = D(I)
140 CONTINUE
    DO 150 I = 1,N
        D(I) = B(I,IC)
        B(I,IC) = B(I,LC)
        B(I,LC) = D(I)
150 CONTINUE
    MC1(IC) = IC
    MC2(IC) = LC
160 IF( B(IC,IC) . EQ . 0.0 ) GO TO 270
    P = B(IC,IC)
    IF ( P . EQ . 1.0 ) GO TO 175
    DO 170 I = 1,M
        B(IC,I) = B(IC,I)/P
170 CONTINUE
175 DO 190 I = 1,N
    IF ( I . EQ . IC ) GO TO 190
    P1 = B(I,IC)
    DO 180 J = IC,M
        B(I,J) = B(I,J) - B(IC,J) * P1

```

```

DO 5 I = 1,N
MC1(I)= 0
MC2(I)= 0
5  CONTINUE
M = N + N
M1 = N + 1
DO 70 I = 1,N
DO 60 J = 1,N
B(I,J) = A(I,J)
60  CONTINUE
70  CONTINUE
DO 90 I = 1,N
DO 80 J = M1,M
B(I,J) = 0.0
80  CONTINUE
90  CONTINUE
DO 100 I = 1,N
M2 = I + N
B(I,M2) = 1.0
100 CONTINUE
IC = 1
110 AL = B(IC,IC)
LR = IC
LC = IC
DO 130 I = IC,N
DO 120 J = IC,N
IF ( ABS (B(I,J) ) . LE . AL ) GO TO 120
AL = B(I,J)

```

```

CY(I,I)=CY(I,I)+CA(I,LJ)+CB(I,LJ)
35  CONTINUE
-25  CONTINUE
    DO 40 I=1,N
    DO 50 LJ=1,N
    IF (I.EQ.LJ) GOTO 50
    CY(I,LJ)=-CA(I,LJ)
50  CONTINUE
40  CONTINUE
    DO 60 I=1,N
    DO 70 LJ=1,N
    Y1(I,LJ)=CABS(CY(I,LJ))
    OO=AIMAG(CY(I,LJ))
    ZZ=REAL(CY(I,LJ))
    IF(OO.EQ.O.O.AND.ZZ.EQ.O.O) GOTO 70
    IF(ZZ.EQ.O.O.AND.OO.NE.O.O) GOTO 4000
    RR=OO/ZZ
    TH(I,LJ)=ATAN(RR)
    GOTO 70
4000 TH(I,LJ) =1.57
70  CONTINUE
60  CONTINUE
    RETURN
    END

```

```

C  ##### MATRIX INVERSE #####
SUBROUTINE MINE(A,C,N)
  DIMENSION A(8,8) , C(8,8) ,      (16)
  DIMENSION MC1(8) , MC2(8) ,      (16)

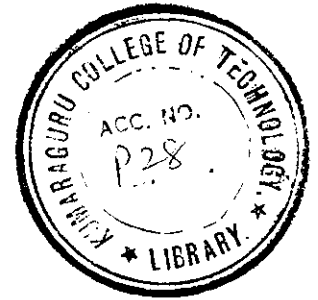
```

```

180  CONTINUE
190  CONTINUE
    IF ( IC . EQ . N ) GO TO 200
    IC = IC + 1
    GO TO 110
200  DO 210 I = 1,N
    DO 210 J = 1,N
        J1 = J + N
        C(I,J) = B(1,J1)
210  CONTINUE
220  CONTINUE
    DO 240 IQ = N,1,-1
        K1=MC1(IQ)
        K2=MC2(IQ)
        IF ( K1 . EQ . 0 ) GO TO 240
        DO 230 J = 1,N
            D(J) = C(K1,J)
            C(K1,J) = C(K2,J)
            C(K2,J) = D(J)
230  CONTINUE
240  CONTINUE
        GO TO 290
270  WRITE (*,280)
280  FORMAT ( 5X,'INVERSE IS NOT EXISTING')
290  CONTINUE
    RETURN
    END

```

C ##### MATRIX MULTIPLICATION #####



INPUT DATA

0502000.00100005
 00.00000
 000.50001.00
 01.0600001.0470001.0381001.0237001.02450
 00.00000-0.04897-0.10740-0.09300-0.03820
 01000.000001000.000001000.000001000.000001000.000001000.000001000.0000
 -1000.0000-1000.0000-1000.0000-1000.0000-1000.0000-1000.0000-1000.0000-1000.0000
 .000 .000 .020 .060 .000 .000 .000 .000 .080 .240
 .020 .060 .000 .000 .040 .120 .060 .180 .060 .180
 .000 .000 .040 .120 .000 .000 .080 .240 .000 .000
 .000 .000 .060 .180 .000 .240 .000 .000 .010 .030
 .080 .240 .060 .180 .000 .000 .010 .030 .000 .000
 .000 .000 .000 .030 .000 .000 .000 .000 .070 .025
 .000 .030 .000 .000 .030 .015 .000 .020 .000 .020
 .000 .000 .000 .015 .000 .000 .000 .025 .000 .000
 .000 .000 .000 .030 .000 .025 .000 .000 .000 .010
 .000 .025 .000 .020 .000 .000 .000 .010 .000 .000
 001.000001.000001.000001.000001.000

```
DO 1 I = 1,L
DO 1 K = 1,N
B(K,I)=A(I,K)
```

```
1 CONTINUE
RETURN
END
```

```
C ##### MATRIX ADDITION #####
```

```
SUBROUTINE ADD(A,B,C,L,N)
DIMENSION A(8,8),B(8,8),C(8,8)
DO 1 I = 1,L
DO 1 K = 1,N
C(I,K) = A(I,K) + B(I,K)
```

```
1 CONTINUE
RETURN
END
```

```
C ##### MATRIX SUBTRACTION #####
```

```
SUBROUTINE SUB (A,B,C,L,N)
DIMENSION A(8,8),B(8,8),C(8,8)
DO 1 I = 1,L
DO 1 K = 1,N
C(I,K) = A(I,K) - B(I,K)
```

```
1 CONTINUE
RETURN
END
```

```
C ##### SUBROUTINES OVER #####
```

```

SUBROUTINE NTNL(A,B,C,L,M,N)
DIMENSION A(B,B) , B(O,O) , C(B,B)
DO 1 I=1,L
DO 2 J=1,N
C(I,J)=O.O
DO 3 K=1,M
C(I,J)=C(I,J)+A(I,K)*B(K,J)
3 CONTINUE
2 CONTINUE
1 CONTINUE
RETURN
END

```

```

C ##### CALCULATING P AND Q #####
SUBROUTINE NR(V,DEL,Y1,TH,N,P,Q)
DIMENSION P(5),Q(5),Y1(10,10),TH(10,10),DEL(5),V(5)
DO 1 I = 1,N
P(I)=O.O
Q(I)=O.O
DO 2 K = 1,N
P(I)=ABS(V(I)*V(K)*Y1(I,K))*COS(TH(I,K)+DEL(I)-DEL(K))+P(I)
Q(I)=-ABS(V(I)*V(K)*Y1(I,K))*SIN(TH(I,K)+DEL(I)-DEL(K))+Q(I)
2 CONTINUE
1 CONTINUE
RETURN
END

```

```

C ##### FINDING TRANSPOSE #####
SUBROUTINE TRANS(A,B,L,N)
DIMENSION A(B,B),B(B,B)

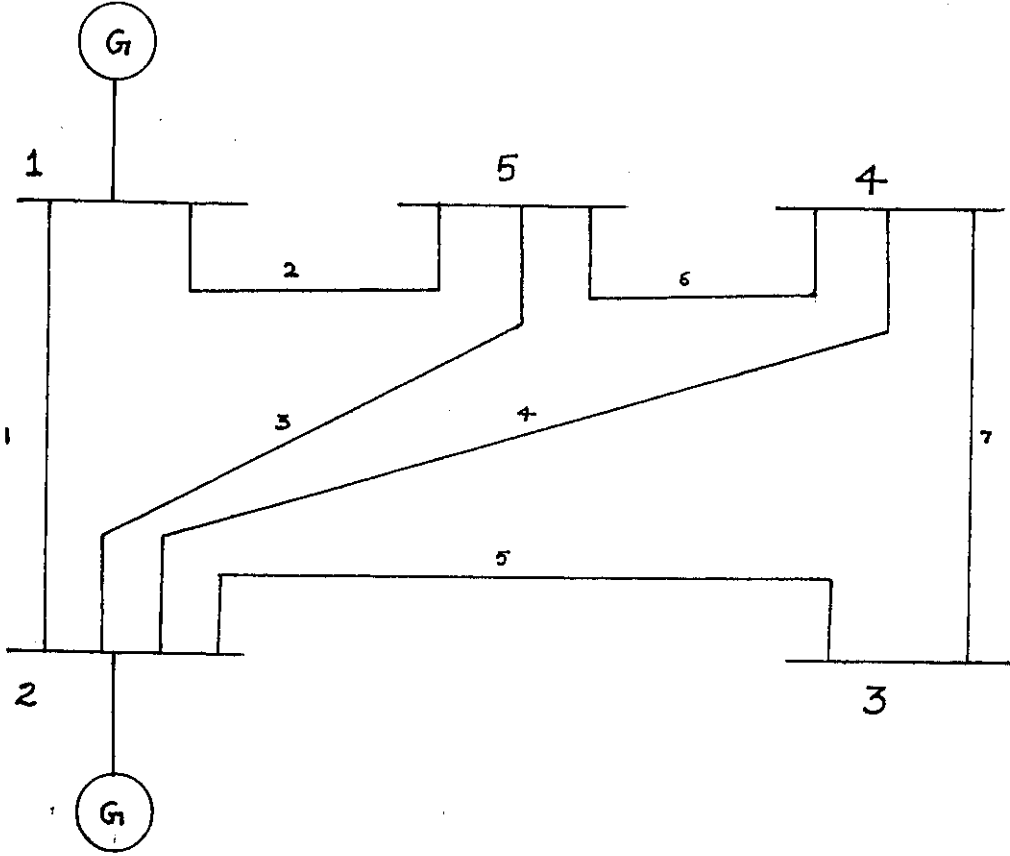
```

RESULTS

VOLTAGE	VOLTAGE ANGLE IN RADIANS	REAL POWER	REACTIVE POWER
1.04600	.00000	-1.02600	.67194
1.04320	-.04590	-.00009	.22250
1.01800	-.10793	.45000	-.46000
1.02530	-.09300	.33950	-.29130
1.02140	-.08610	.30490	-.39153

formula used

FIVE BUS TEST SYSTEM



BUS DATA

Bus Number	Bus Code	Initial Bus voltage		Generation		Load	
		Mag P.U	Angle P.U	MW	MUAR	MW	MUAR
1*	1	1.060	0	0	0	0	0
2	2	1.047	0	40	30	20	10
3	3	1.000	0	0	0	60	10
4	4	1.000	0	0	0	40	5
5	5	1.000	0	0	0	45	15

conclusion

CONCLUSION

A new concept involving an optimal approach to the use of load flow as a planning tool is presented in this work. Normally at the planning stage, the system planner has to carry out a number of successive load flow studies, each time varying the specification, before an acceptable operating condition is attained. The success of this procedure depends almost entirely on the skill and experience of the planner and his knowledge of the system. The new approach suggested here eliminates this drawback and uses an optimisation technique to arrive at the best solution in a single run of the program.

The method was implemented using a Fortran program on a sample 5 bus (thee IEEE standard) system. Convergence was obtained in 3 iteration to a tolerance of 0.001 per unit. The results were very encouraging and the authors feel reasonably certain that the method can be applied with success to large systems also.

The main aim of the present work was to illustrate the validity of the new approach. The computations effort required is more than for a conventional load flow and further work needs to be carried out to improve the efficiency and speed of the program.

bibliography

BIBLIOGRAPHY

1. "COMPUTER METHODS IN POWER SYSTEM ANALYSIS"
BY STAGG AND EI _ ABIAD.
2. "ELECTRICAL POWER SYSTEMS"
BY C.L. WADHWA
3. "JOURNAL OF THE INTSITUTION OF ENGINEERS (India)"
VOLUME 69, AUGUST 1 988