



**An Optimized Polymorphic Hybrid Multicast  
Routing Protocol for MANET's**



*P - 2171*

***A PROJECT REPORT***

*Submitted by*

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**ANNA UNIVERSITY: CHENNAI 600 025**

***APRIL 2008***

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## BONAFIDE CERTIFICATE

Certified that this project report “An Optimized Polymorphic Hybrid Multicast Routing Protocol for MANET’s” is the bonafide work of “Revathi R and Saraswathi P” who carried out the project work under my supervision.



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## DECLARATION

We,

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hereby declare that the project entitled “An Optimized Polymorphic Hybrid Multicast Routing Protocol for MANET’s”, submitted in partial fulfillment to Anna University as the project work of Bachelor of Technology (Information Technology) degree, is a record of original work done by us under the supervision and guidance of Department of Information Technology, Kumaraguru college of Technology, Coimbatore.

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## ABSTRACT

We propose in this paper an optimized, polymorphic, hybrid multicast routing protocol for MANET. This new polymorphic protocol attempts to benefit from the high efficiency of proactive behavior (in terms of quicker response to transmission requests) and the limited network traffic overhead of the reactive behavior, while being power, mobility, and vicinity-density (in terms of number of neighbor nodes per specified area around a mobile node) aware. The proposed protocol is based on the principle of adaptability and multibehavioral modes of operations. It is able to change behavior in different situations in order to improve certain metrics like maximizing battery life, reducing communication delays, improving deliverability, etc. The protocol is augmented by an optimization scheme, adapted from the one proposed for the Optimized Link State Routing protocol (OLSR) in which only selected neighbor nodes propagate control packets to reduce the amount of control overhead. Extensive simulations and comparison to peer protocols demonstrated the effectiveness of the proposed protocol in improving performance and in extending battery power longevity.

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## LIST OF ABBREVIATIONS

- OPHMR** - An Optimized Polymorphic Hybrid  
Multicast Routing Protocol for MANET's
- GLOMOSIM** - GLObal MObile Information System  
SIMulator
- PARSEC** - PARallel Simulation Environment for  
Complex systems
- MPR** - MultiPoint Relay
- MANET** - Mobile Ad-hoc NETWORK

# 1.INTRODUCTION

## 1.1 GENERAL

Mobile ad hoc networks (MANETs) are infrastructure free networks of mobile nodes that communicate with each other wirelessly. There are several routing schemes that have been proposed and several of these have been extensively simulated or completely implemented as well. The primary applications of such networks have been in disaster relief operations, military use, conferencing and environment sensing. Unlike conventional wireless networks one may find in offices, universities, communities or homes there is no central entity that controls how, when and where, packets are delivered to each recipient. All communication takes place in an ad hoc manner, which means on the fly and all the nodes in the network participate in relaying packets or messages to each other whenever it is possible for each node to do so.

There are several unicast routing algorithms that have been developed for MANETs that have their own unique characteristic strengths and weaknesses. Most protocols can be classified in several ways. Some are classified as reactive or on- demand while others are proactive. In general, a proactive protocol finds routes in advance while a reactive protocol finds routes to the destination only when it absolutely must. For example, Ad hoc On demand Distance Vector routing (AODV) [AODV] is an on-demand protocol since no protocol information is transmitted before an application decides to send data and no data is sent until a route is formed, whereas Destination Sequenced Distance Vector protocol (DSDV) [DSDV] is a more proactive protocol in which routes are discovered and stored even before they are needed.

Proactive protocols generally generate much more traffic than on-demand protocols. A third general category is a hybrid algorithm that effectively combines multiple characteristics in a unique and meaningful way. For example, the Zone Routing Protocol (ZRP) [ZRP] is a hybrid protocol that combines local proactive routing with a globally reactive routing strategy.

## 1.2 PROBLEM DEFINITION

This new polymorphic protocol attempts to benefit from the high efficiency of proactive behavior (in terms of quicker response to transmission requests) and the limited network traffic overhead of the reactive behavior, while being power, mobility, and vicinity-density (in terms of number of neighbor nodes per specified area around a mobile node) aware. The proposed protocol is based on the principle of adaptability and multibehavioral modes of operations. It is able to change behavior in different situations in order to improve certain metrics like maximizing battery life, reducing communication delays, improving deliverability, etc. The protocol is augmented by an optimization scheme, adapted from the one proposed for the Optimized Link State Routing protocol (OLSR) in which only selected neighbor nodes propagate control packets to reduce the amount of control overhead. Extensive simulations and comparison to peer protocols demonstrated the effectiveness of the proposed protocol in improving performance and in extending battery power longevity.

### 1.3 OBJECTIVE OF THE PROJECT

The simple hybrid routing protocol is built on the principle that the mobile node is able to behave either proactively or reactively under different conditions. As such, the resulting protocol inherits the benefits of both reactive and proactive behaviors and, hence, can expand its suitability region within the ad hoc network design space.

The distance of the destination nodes from the source nodes are calculated and based on that, the mode of transmission (either proactive or reactive ) is chosen. In case of short distance transmission ,the reactive mode is used and in long distance transmission proactive mode is used.

In the OPHMR protocol, each node maintains a two-hop Neighborhood Table (2NTable). The 2NTable is used to calculate the MPR information. When a node receives an update packet, it uses the neighborhood information in the packet to calculate the two-hop neighborhood and updates the corresponding entries in the 2NTable.

MPR nodes are selected to forward broadcast messages during the flooding process. This technique substantially reduces the message overhead as compared to the classical flooding mechanism, where every node retransmits each message when it receives it for the first time. Each node has its MPR set and will broadcast its MPR information in the periodic update packets. When propagating the periodic update packets, only the MPRs forward update packets.

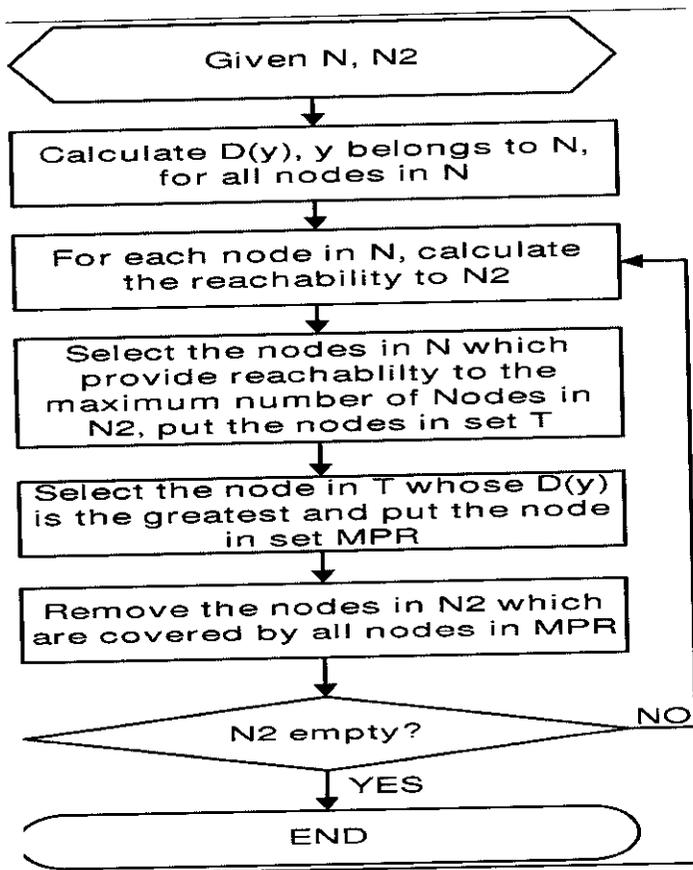


Fig 1.3.1 Flowchart for the MPR operation.

## 2. LITERATURE REVIEW

### 2.1 FEASIBILITY STUDY

#### 2.1.1 CURRENT STATUS OF THE PROBLEM

Proactive protocols continuously exchange network topology information so as to constantly monitor topology changes and use that knowledge for efficient, low latency data transmission. While in reactive protocols nodes react only on-demand to data transmission requests and perform path finding operations only when needed.

##### **Limitations**

The network topology of ad hoc networks is unstable and changes frequently with nodes' mobility, traditional routing protocols in static networks are not efficient for ad hoc networks.

#### 2.1.2 PROPOSED SYSTEM

The simple hybrid routing protocol is built on the principle that the mobile node is able to behave either proactively or reactively under different conditions. As such, the resulting protocol inherits the benefits of both reactive and proactive behaviors and, hence, can expand its suitability region within the ad hoc network design space.

OPHMR is built using the proactive behavior of the MZR protocol and the reactive behavior of ODMRP. In addition, the protocol was augmented with the Multipoint Relay (MPR)-based optimization mechanism of OLSR [4]. This is done with the objective of reducing the number of control packets forwarded. Next, we describe its proactive and reactive behaviors as well as the adopted optimized forwarding mechanism.

**PROACTIVE BEHAVIOR.** When a node is in PM1 or PM2, it periodically sends out update packets which have a TTL set to the Zone Radius. When a node receives an update packet, if it is in PM1, in PM2, or in PRM, it saves the information in the packet into the one-hop Neighborhood Table (NTable), reduces the TTL by 1, and forwards the packet.

**REACTIVE BEHAVIOR.** When a node has packets to send to a multicast group or wants to join a multicast group, it sends out a Join\_Request packet and waits for its Join\_Reply from a destination node. Only nodes that belong to the destination multicast group will send out Join\_Replies, and such nodes will update their Multicast Routing Tables (MRTable) to maintain the multicast group information. If an intermediate node (including the source node) has entries in its NTable that belong to the destination multicast group, it unicasts the join\_Request to the nodes registered in those entries. When the source node receives a Join\_Reply, it updates its MRTable and begins data transmission.

**ROUTING TABLES:** Each node maintains the two routing tables, the NTable and the MRTable. The NTable acts as the neighborhood routing table we described in the algorithm and, actually, only nodes in Proactive Modes maintain it. Each neighbor in the zone has an entry in this table. Each such an entry contains the routing information to the node, including hop count and next hop address. In addition, it contains the multicast routing information of the node, such as the multicast group that the node belongs to. Each entry is assigned a lifetime and the one that exceeds its predefined lifetime is removed from the NTable. Nodes in a PM or a PRM mode maintain the NTable using the update packets they receive from others. Nodes in those modes also periodically flush the NTable to remove stale entries. Nodes in RM mode just periodically flush their NTables to remove

stale entries, but will not update or insert any entries in them. Each node also uses its MRTable to maintain both its multicast routing information and the multicast routing topology. The structure of the MRTable is the same as the routing table used for ODMRP [32].

**PACKET STRUCTURE.** The structure of zone UPDATE packets used in the OPHMR protocol is shown in Fig. 1. In this table, the type is set to 4 to indicate that the packet is for ZONE UPDATE. The field “Reserved” is used for future use. The field “Time To Live” has the same classical usage. When the source node generates a ZONE UPDATE packet, it sets its Time To Live to the Zone\_Radius. Its value is decremented by one at each hop until it becomes zero and it is then discarded. The field “Number of Multicast Groups” indicates how many multicast groups this node belongs to. The field “Source Node IP Address” stores the IP address of the source node. Finally, each “Multicast Group IP Address” field describes information on a particular multicast group that the source node belongs to. The second type of packet we describe here is the notification packet. Its structure is very simple. The first eight bits are for the packet type and are set to 5 to indicate that it is a Notification Packet for OPHMR. The next 16 bits are for the switch type and each bit is for one type (there are 12 types of switching, so the first four bits are always set to 0). The last eight bits are reserved. Table 1 shows the meaning of each bit.

0										1										2										3			
0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1		
Type										Reserved										Time To Live										Number of Multicast Groups			
Source Node IP Address																																	
Multicast Group IP Address (1)																																	
Multicast Group IP Address (2)																																	
.....																																	

Fig 2.1.2.1: Packet Structure of OPHMR

**Advantages:**

The protocol attempts to benefit from the high efficiency of proactive routing in reducing response time to transmission requests and from the reduced control overhead offered by reactive routing.

**2.2 HARDWARE SPECIFICATION**

- Processor : Intel Pentium IV 2.2 GHZ
- Hard Disk Drive : 20 GB.
- RAM : 256 MB RAM minimum.

**2.3 SOFTWARE REQUIREMENTS**

- Operating System : Windows 2000/XP
- Software : Glomosim Simulator

## **2.4 SOFTWARE OVERVIEW**

### **2.4.1 GLOMOSIM**

Our project has been done using a scalable simulation environment called GloMoSim (for Global Mobile Information System Simulator) that effectively utilizes parallel execution to reduce the simulation time of detailed high-fidelity models of large communication networks. GloMoSim has been designed to be extensible and composable: the communication protocol stack for wireless networks is divided into a set of layers, each with its own API. Models of protocols at one layer interact with those at a lower (or higher) layer only via these APIs. The modular implementation enables consistent comparison of multiple protocols at a given layer. The parallel implementation of GloMoSim can be executed using a variety of conservative synchronization protocols, which include the null message and conditional event algorithms.

### **2.4.2 PARSEC**

PARSEC (for PARAllel Simulation Environment for Complex systems) is a C-based simulation language developed by the Parallel Computing Laboratory at UCLA, for sequential and parallel execution of discrete-event simulation models. It can also be used as a parallel programming language. PARSEC runs on several platforms, including most recent UNIX variants as well as Windows. PARSEC adopts the process interaction approach to discrete-event simulation. An object (also referred to as a physical process) or set of objects in the physical system is represented by a logical process. Interactions among physical processes (events) are modeled by time-stamped message exchanges among the corresponding logical processes. One of the important distinguishing features of PARSEC

is its ability to execute a discrete-event simulation model using several different asynchronous parallel simulation protocols on a variety of parallel architectures. PARSEC is designed to cleanly separate the description of a simulation model from the underlying simulation protocol, sequential or parallel, used to execute it. Thus, with few modifications, a PARSEC program may be executed using the traditional sequential (Global Event List) simulation protocol or one of many parallel optimistic or conservative protocols. In addition, PARSEC provides powerful message receiving constructs that result in shorter and more natural simulation programs.

### **2.4.3 ABOUT GLOMOSIM**

GloMoSim is a mobile simulator built using C language. All message transfers and other network elements are handled by the individual layer coding built in C. To make the concepts clear, GloMoSim provides users with a Visualization Tool ( VT ) built using Java. The VT helps us understand the network environment, the node positions, message transfers, clustering details, etc.

### **2.4.4 GLOMOSIM ARCHITECTURE**

The networking stack is decomposed into a number of layers as shown in Figure 1. A number of protocols have been developed at each layer and models of these protocols or layers can be developed at different levels of granularity.

In our project we deal with all these layers, but most of the coding has been implemented in the Clustering and Routing layers. The dynamic clustering algorithm explained in chapters 2 and 3 is implemented in these two layers. The clustering algorithm has been fully implemented in the

Clustering layer. The coding present in the Routing layer has been largely modified according to the cluster formation. The algorithm used for routing is the Bellman-Ford algorithm. This algorithm maintains a routing table for every node being simulated. We have modified the entries of this table based on the cluster formation which is dynamic.

The communication between the various layers is accomplished by means of the various APIs available. A common API between two layers helps in the communication between those two layers.

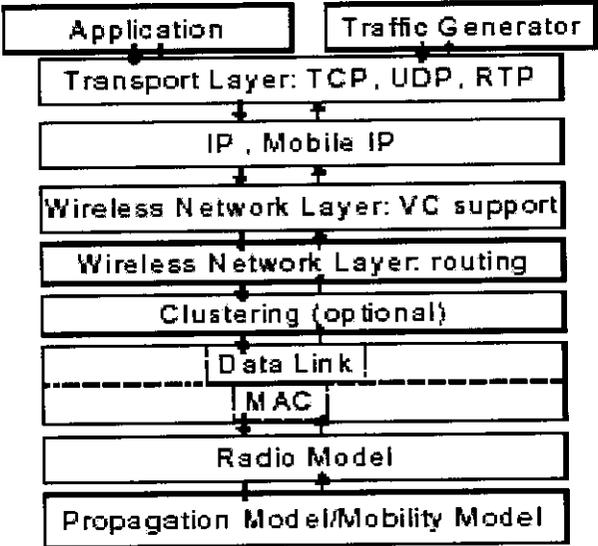


Fig.2.4.4.1. GloMoSim Architecture



## 2.4.5 GLOMOSIM LIBRARY

GloMoSim is a scalable simulation library for wireless network systems built using the PARSEC simulation environment. Table 1 lists the GloMoSim models currently available at each of the major layers. GloMoSim also supports two different node mobility models. Nodes can move according to a model that is generally referred to as the “random waypoint” model. A node chooses a random destination within the simulated terrain and moves to that location based on the speed specified in the configuration file. After reaching its destination, the node pauses for a duration that is also specified in the configuration file. The other mobility model in GloMoSim is referred to as the “random drunken” model. A node periodically moves to a position chosen randomly from its immediate neighboring positions. The frequency of the change in node position is based on a parameter specified in the configuration file.

Layer:	Models:
Physical (Radio propagation)	Free space, Rayleigh, Ricean, SIRCIM
Data Link (MAC)	CSMA, MACA, MACAW, FAMA, 802.11
Network (Routing)	Flooding, Bellman-Ford, OSPF, DSR, WRP
Transport	TCP, UDP
Application	Telnet, FTP

**Table 2.4.5.1: Models currently in the GloMoSim library.**

### 3. DETAILS OF METHODOLOGY EMPLOYED

#### 3.1 ALGORITHMS

##### Algorithm 1

##### POLYMORPHIC ALGORITHM

```
if Power > P TH1 then
if the MN is not in PM1, it switches to PM1.
then it notifies neighbors about the mode switch.
else
if Power < P TH2 then
if the MN is not in RM, it switches to RM.
then it notifies neighbors about the mode switch.
else
Perform the mobility speed routine.
end if
end if
```

##### Algorithm 2

##### MOBILITY SPEED ROUTINE

```
if Mobility > M TH then
if Vicinity < V TH then
if the MN is not in PM2, it switches to PM2.
then it notifies neighbors about the mode switch.
else
if the MN is not in PRM switches to PRM.
then it notifies neighbors about the mode switch.
```

end if  
else  
if the MN is not in RM switches to RM.  
then it notifies neighbors about the mode switch.  
end if

### 3.2. Path Finding Procedure

When a node has packets to send to a multicast group or wants to join that group, it begins the path finding procedure. If it is in RM mode, it sends out a Join\_Request the way it is done in ODMRP and waits for replies. If the node is in proactive mode or proactive ready mode, it first looks in its NTable to see whether there are nodes that belong to the destination multicast group. If so, it unicasts Join\_Requests to all these nodes and waits for replies. Otherwise, the node will broadcast a Join\_Request. When a node receives a Join\_Request and it is a member of the multicast group, it generates a reply and sends it back to the source of the Join\_Request, updating the MRTable to record the route. If the node could not send a reply, it checks its own behavior. If it is in a reactive mode, it just propagates the Join\_Request and records it in the route cache. If the node is in a proactive mode or in the PRM, it looks in its own NTable to find the destination multicast group member. If there are members in its zone, it unicasts the Join\_Request to all of them. If not, it just propagates the Join\_Request. When the source node receives a reply, it updates its MRTable to record the route and begins data transmission.

## 4. PERFORMANCE EVALUATION

### 4.1 SIMULATION ENVIRONMENT

An extensive simulations are conducted to evaluate the performance of OPHMR. The simulations were implemented on GloMoSim 2.03, a scalable simulation environment for wireless network systems. GloMoSim uses the parallel discrete-event simulation capability provided by PARSEC.

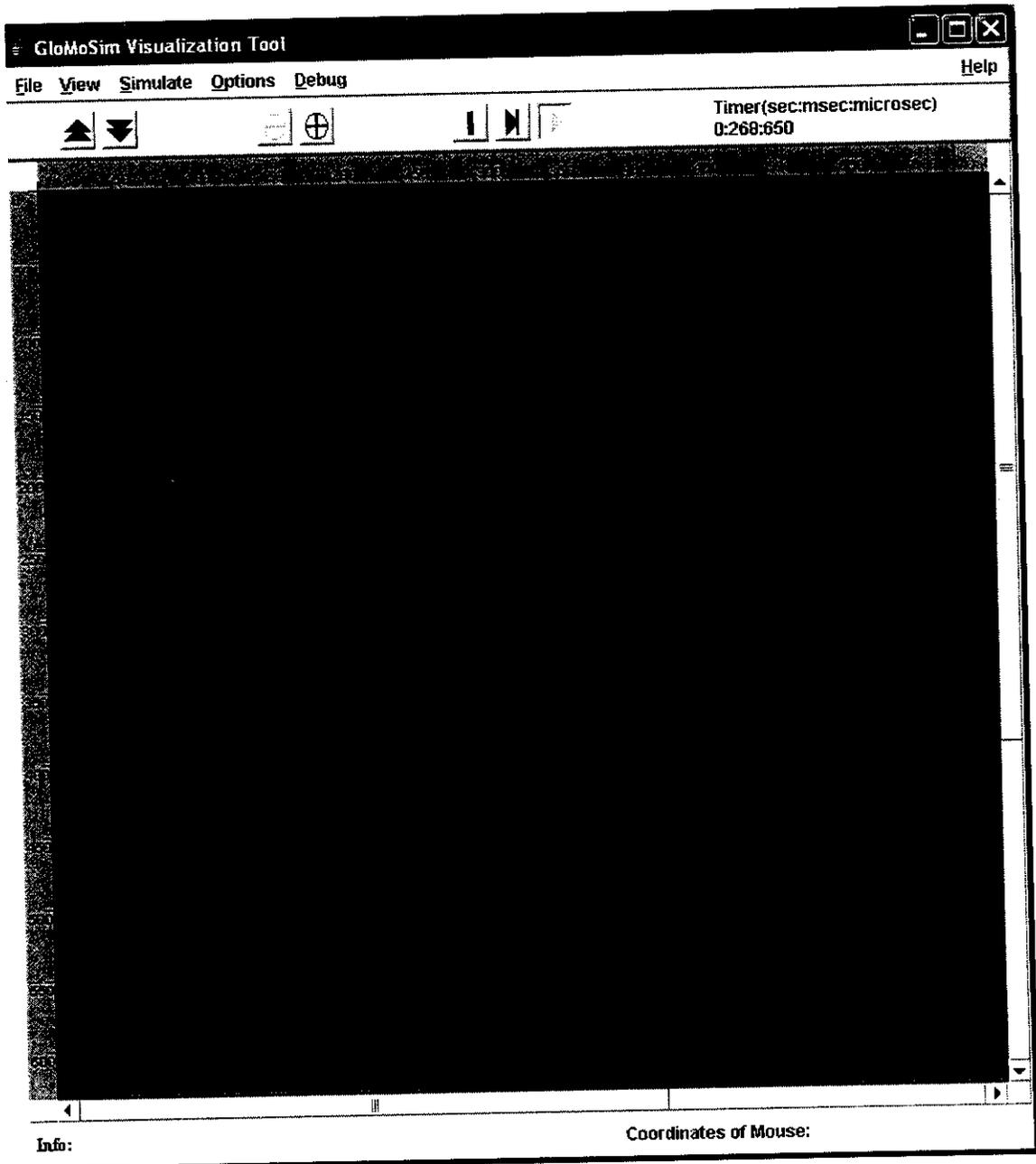
The simulation parameters used are as follows:

PARAMETERS	VALUES
Network size	49nodes
Radio Range	250m
Packet size	512bytes
Terrain-Dimensions	(500,500)
Node-Placement	Uniform
GUI Option	YES
Routing Protocols	OPHMR

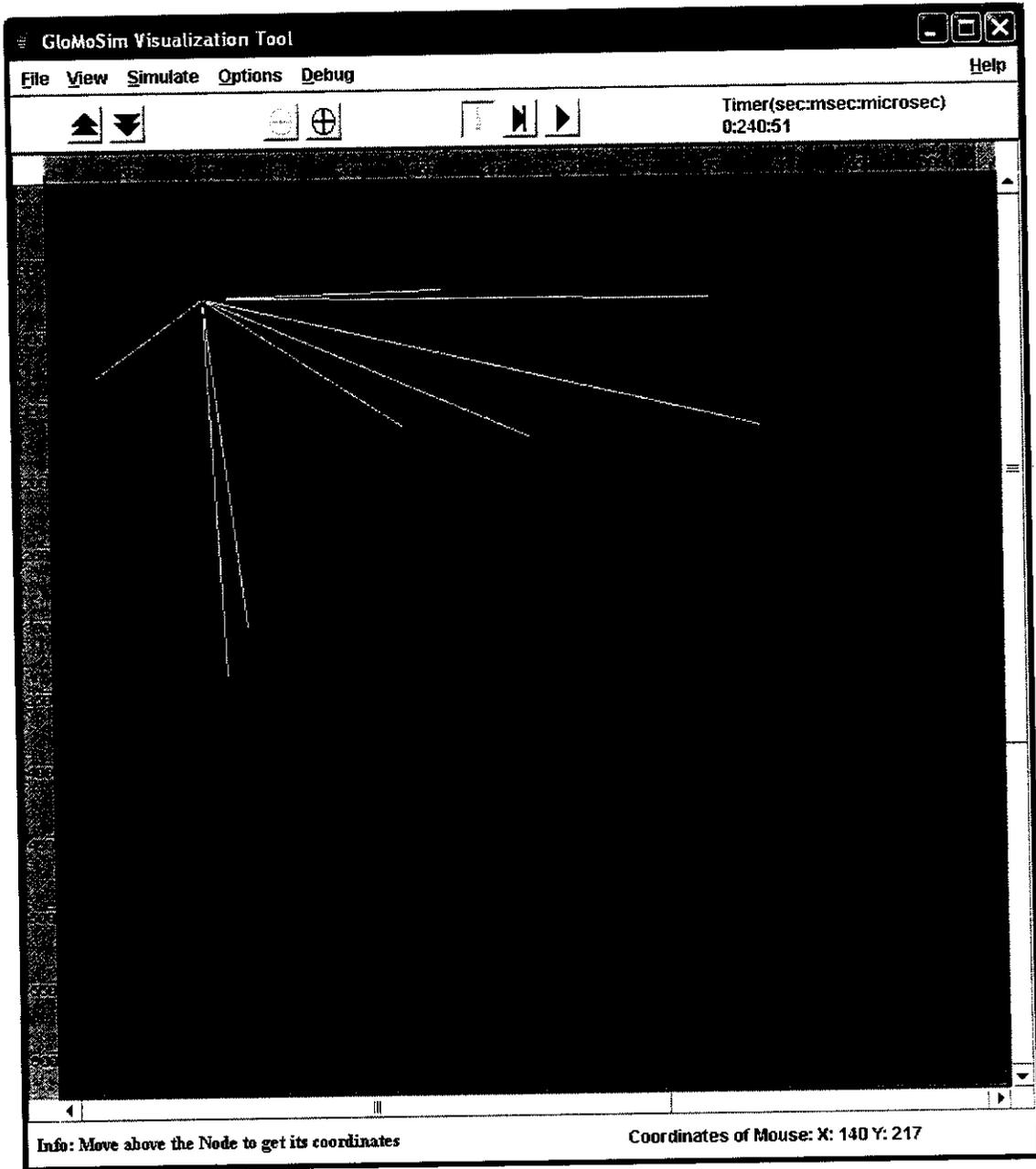
Table.4.1.1 Simulation Parameters

# SCREEN SHOTS

## 1. NODE PLACEMENT

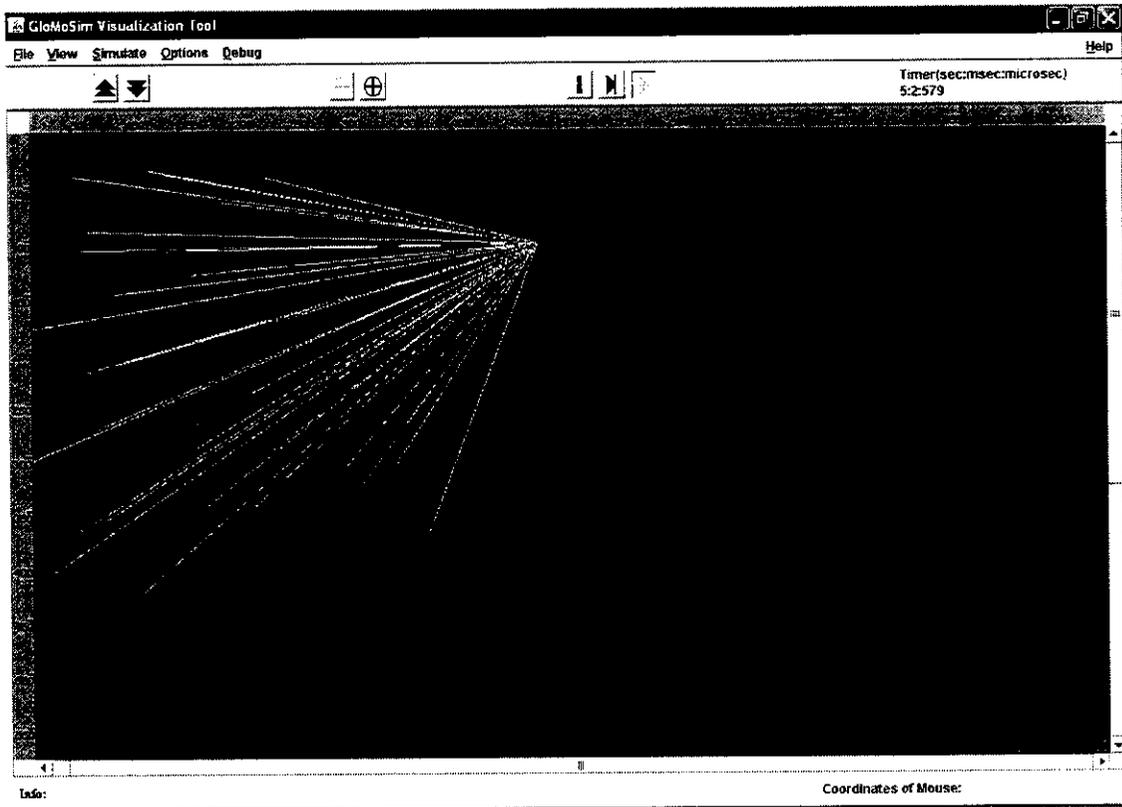


## 2. CONNECTION ESTABLISHMENT

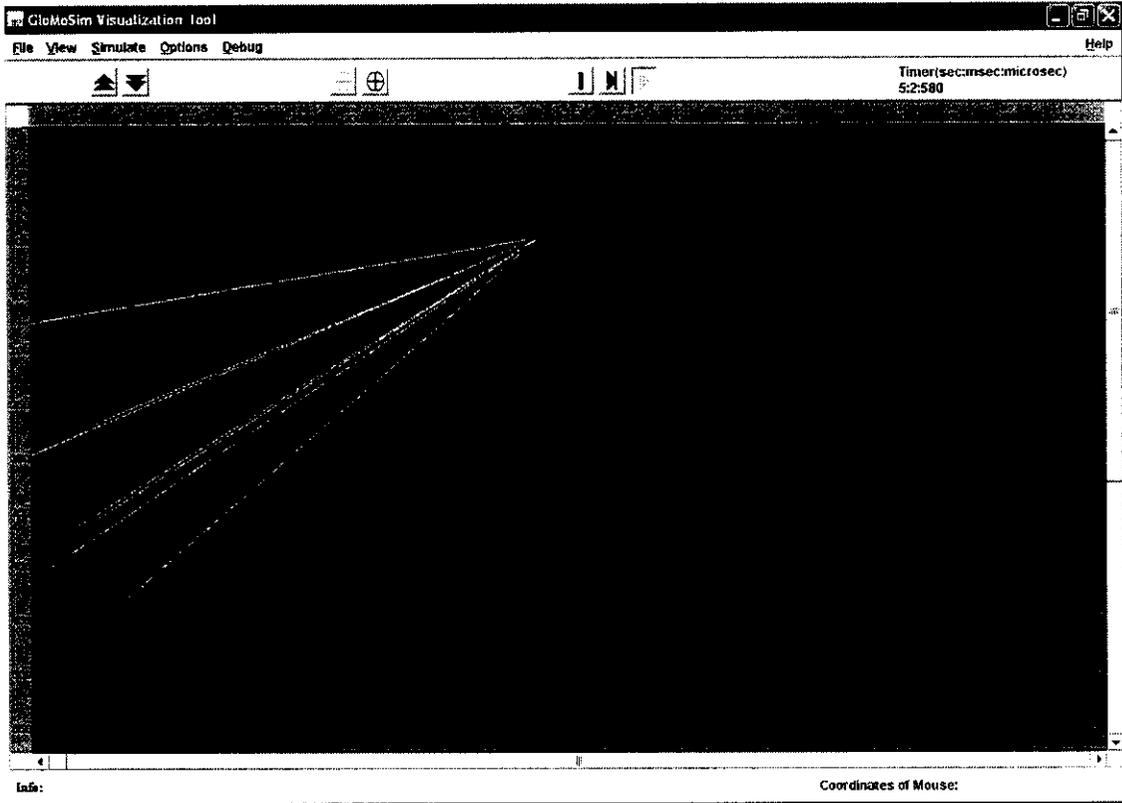


## 4.2 SIMULATION RESULTS

### SHORT DISTANCE TRANSMISSION(REACTIVE)



# LONG DISTANCE TRANSMISSION(PROACTIVE)





## 5.CONCLUSION

The protocol design is a novel way of combining three dimensions in protocol design, namely, hybridity, adaptability, and power awareness. With regard to hybridity, The protocol attempts to benefit from the high efficiency of proactive routing in reducing response time to transmission requests and from the reduced control overhead offered by reactive routing.

## 6.FUTURE ENHANCEMENT

The project has been further enhanced to include the mechanisms for the battery power calculation of every nodes and can change the mode of transmission .This new concept of polymorphic protocols constitutes the next trend in the design of efficient multibehavioral routing protocols for wireless, power-constrained networks such as MANETs.

## APPENDICES

### SOURCE CODE:

```
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#include <assert.h>
#include <math.h>
#include "api.h"
#include "structmsg.h"
#include "fileio.h"
#include "main.h"
#include "message.h"
#include "network.h"
#include "ip.h"
#include "nwip.h"
#include "nwcommon.h"
#include "application.h"
#include "transport.h"
#include "java_gui.h"

#include "mac.h"
#include "ophmr.h"

#define noDEBUG

void RoutingOphmrInit(GlomoNode *node,
                     GlomoRoutingOphmr **ophmrPtr,
                     const GlomoNodeInput *nodeInput)
{
    GlomoNodeInput memberInput;
    BOOL retVal;
    int i;

    char joinTimeStr[GLOMO_MAX_STRING_LENGTH];
    char leaveTimeStr[GLOMO_MAX_STRING_LENGTH];
    clocktype joinTime, leaveTime;
    NODE_ADDR srcAddr, mcastAddr;

    GlomoRoutingOphmr *ophmr =
```

```

    (GlomoRoutingOphmr *)checked_pc_malloc
(sizeof(GlomoRoutingOphmr));

(*ophmrPtr) = ophmr;

if (ophmr == NULL)
{
    fprintf(stderr, "OPHMR: Cannot alloc memory for OPHMR
    struct!\n");
    assert (FALSE);
}

NetworkIpSetRouterFunction(node, &RoutingOphmrRouterFunction);

RoutingOphmrInitStats(node);
RoutingOphmrInitMembership(&ophmr->memberFlag);
RoutingOphmrInitFgFlag(&ophmr->fgFlag);
RoutingOphmrInitMemberTable(&ophmr->memberTable);
RoutingOphmrInitTempTable(&ophmr->tempTable);
RoutingOphmrInitRouteTable(&ophmr->routeTable);
RoutingOphmrInitMessageCache(&ophmr->messageCache);
RoutingOphmrInitSeqTable(node);
RoutingOphmrInitSent(&ophmr->sentTable);
RoutingOphmrInitAckTable(&ophmr->ackTable);
RoutingOphmrInitResponseTable(&ophmr->responseTable);

retVal = GLOMO_ReadCachedFile(nodeInput, "MCAST-CONFIG-
FILE",
    &memberInput);

if (retVal == FALSE)
{
    fprintf(stderr, "OPHMR: Needs MCAST-CONFIG-FILE.\n");
    assert(FALSE);
}

for (i = 0; i < memberInput.numLines; i++)
{
    retVal = sscanf(memberInput.inputStrings[i],
        "%ld %ld %s %s",&srcAddr, &mcastAddr,

```

```
joinTimeStr, leaveTimeStr);
```

```
if (retVal != 4)
```

```
{  
    fprintf(stderr, "Application: Wrong configuration  
    format!\n");  
    assert(0); abort();  
}
```

```
joinTime = GLOMO_ConvertToClock(joinTimeStr);  
leaveTime = GLOMO_ConvertToClock(leaveTimeStr);
```

```
if (node->nodeAddr == srcAddr)
```

```
{  
    RoutingOphmrSetTimer( node, MSG_NETWORK_JoinGroup,  
    mcastAddr, joinTime);  
    RoutingOphmrSetTimer(  
    node, MSG_NETWORK_LeaveGroup, mcastAddr, leaveTime);  
}
```

```
}  
void RoutingOphmrHandleProtocolEvent(GlomoNode *node, Message  
*msg)
```

```
{  
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node->  
    networkData.networkVar;  
    GlomoRoutingOphmr* ophmr= (GlomoRoutingOphmr *) ipLayer->  
    routingProtocol;  
    char clockStr[GLOMO_MAX_STRING_LENGTH];
```

```
switch (msg->eventType)
```

```
{  
    case MSG_NETWORK_CheckAcked:
```

```
{  
    NODE_ADDR *mcastAddr = (NODE_ADDR  
    *)GLOMO_MsgReturnInfo(msg);  
    clocktype jitterTime;
```

```
    if (RoutingOphmrCheckAckTable(*mcastAddr, &ophmr->  
    ackTable))
```

```

{
    #ifdef DEBUG
    printf("Node %ld not acked. Retx!\n", node->nodeAddr);
    #endif
    RoutingOphmrRetxReply(node,*mcastAddr,&ophmr->ackTable);
}

GLOMO_MsgFree(node, msg);

break;
}

case MSG_NETWORK_CheckTimeoutAlarm:
{
    NODE_ADDR *mcastAddr=(NODE_ADDR
        *)GLOMO_MsgReturnInfo(msg);
    if(RoutinOphmrLookupSentTable(*mcastAddr,&ophmr->sentTable))
    {
        if (RoutingOphmrCheckMinExpTime(
            *mcastAddr, &ophmr->sentTable) &&
            RoutingOphmrCheckCongestionTime(
            *mcastAddr, &ophmr->memberTable) &&
            !RoutingOphmrCheckSendQuery(*mcastAddr,&ophmr->
                sentTable))
        {
            #ifdef DEBUG
            printf("HERE\n");
            #endif
            RoutingOphmrSetSendQuery(*mcastAddr,&ophmr->
                sentTable);
        }

        RoutingOphmrSetTimer(node,
            MSG_NETWORK_CheckTimeoutAlarm,
            *mcastAddr, OPHMR_TIMER_INTERVAL);
    }
    GLOMO_MsgFree(node, msg);
    break;
}
case MSG_NETWORK_JoinGroup:

```

```

{
    NODE_ADDR *mcastAddr = (NODE_ADDR*)
        GLOMO_MsgReturnInfo(msg);
    RoutingOphmrJoinGroup(node, *mcastAddr);
    GLOMO_MsgFree(node, msg);
    break;
}

```

```

case MSG_NETWORK_LeaveGroup:
{
    NODE_ADDR *mcastAddr = (NODE_ADDR*)
        GLOMO_MsgReturnInfo(msg);
    RoutingOphmrLeaveGroup(node, *mcastAddr);
    GLOMO_MsgFree(node, msg);
    break;
}

```

```

case MSG_NETWORK_SendReply:
{
    NODE_ADDR *mcastAddr = (NODE_ADDR*)
        GLOMO_MsgReturnInfo(msg);
    RoutingOphmrSendReply(node, *mcastAddr, &ophmr-
        >memberTable, &ophmr->tempTable);
    GLOMO_MsgFree(node, msg);
    break;
}

```

```

case MSG_NETWORK_CheckFg:
{
    NODE_ADDR *mcastAddr = (NODE_ADDR*)
        GLOMO_MsgReturnInfo(msg);
    if (RoutingOphmrLookupFgFlag(*mcastAddr, &ophmr->fgFlag))
    {
        if (RoutingOphmrCheckFgExpired(*mcastAddr, &ophmr->fgFlag))
        {
            RoutingOphmrResetFgFlag(*mcastAddr, &ophmr->fgFlag);
        }
    }
    else
    {

```

```

        RoutingOphmrSetTimer(node, MSG_NETWORK_CheckFg,
            *mcastAddr, OPHMR_FG_TIMEOUT);
    }
}
GLOMO_MsgFree(node, msg);
break;
}

default:
    printf("Time %s: Node %ld received message of unknown type
        %d.\n", clockStr, node->nodeAddr, msg->eventType);
    assert(FALSE);
}
}

void RoutingOphmrHandleProtocolPacket(GlomoNode *node, Message
    *msg, NODE_ADDR srcAddr, NODE_ADDR destAddr)
{
    OPHMR_PacketType *ophmrHeader = (OPHMR_PacketType *)
        GLOMO_MsgReturnPacket(msg);

    switch (*ophmrHeader)
    {
        case OPHMR_ACK:
            RoutingOphmrHandleAck(node, msg, srcAddr, destAddr);
            break;

        case OPHMR_JOIN_REPLY:
            RoutingOphmrHandleReply(node, msg, srcAddr, destAddr);
            break;

        default:
            printf("OPHMR received packet of unknown type\n");
            assert (FALSE);
    }
}

void RoutingOphmrRouterFunction(GlomoNode *node, Message
    *msg, NODE_ADDR destAddr, BOOL
    *packetWasRouted)

```

```

{
GlomoNetworkIp* ipLayer = (GlomoNetworkIp *)node->
                        networkData.networkVar;
GlomoRoutingOphmr* ophmr =(GlomoRoutingOphmr *)ipLayer-
                        >routingProtocol;
IpHeaderType *ipHeader = (IpHeaderType *) msg->packet;
OphmrIpOptionType option;

if (ipHeader->ip_p == IPPROTO_OPHMR)
{
    if (ipHeader->ip_dst >= IP_MIN_MULTICAST_ADDRESS)
    {
        NODE_ADDR sourceAddress;
        NODE_ADDR destinationAddress;
        unsigned char IpProtocol;
        unsigned int ttl;
        NetworkQueueingPriorityType priority;

        NetworkIpRemoveIpHeader(node, msg, &sourceAddress,
            &destinationAddress, &priority, &IpProtocol, &ttl);
        RoutingOphmrHandleProtocolPacket(
            sourceAddress,destinationAddress);
        *packetWasRouted = TRUE;
        return;
    }
    else
    {
        *packetWasRouted = FALSE;
        return;
    }
}

*packetWasRouted = TRUE;
if (FindAnIpOptionField(ipHeader, IPOPT_OPHMR) == NULL &&
    ipHeader->ip_src == node->nodeAddr)
{

#ifdef DEBUG
printf("Node %ld has data to send\n", node->nodeAddr);
#endif
}

```

```

if (RoutingOphmrLookupMembership(destAddr, &ophmr-
>memberFlag))
{
if (RoutingOphmrLookupSentTable(destAddr, &ophmr->sentTable)
&&
!RoutingOphmrCheckSendQuery(destAddr, &ophmr->sentTable))
{
#ifdef DEBUG
printf(" Sending Data\n ");
#endif
RoutingOphmrSendData(node, msg, destAddr);
}
else
{
#ifdef DEBUG
printf(" Sending Join Query\n");
#endif
RoutingOphmrSendQuery(node, msg, destAddr);
}
}
}
else
{
RoutingOphmrHandleData(node, msg);
}
}

```

```

void RoutingOphmrFinalize(GlomoNode *node)

```

```

{
GlomoNetworkIp *ipLayer = (GlomoNetworkIp *)node-
>networkData.networkVar;
GlomoRoutingOphmr *ophmr = (GlomoRoutingOphmr *)ipLayer-
>routingProtocol;

FILE *statOut;
char buf[GLOMO_MAX_STRING_LENGTH];

sprintf(buf, "Number of Join Queries Txed = %d",
ophmr->stats.numQueryTxed);
GLOMO_PrintStat(node, "RoutingOphmr", buf);
sprintf(buf, "Number of Join Replies Txed = %d",

```

```

        ophmr->stats.numReplySent);
GLOMO_PrintStat(node, "RoutingOphmr", buf);
sprintf(buf, "Number of Acks Txed = %d",
        ophmr->stats.numAckSent);
GLOMO_PrintStat(node, "RoutingOphmr", buf);

sprintf(buf, "Number of CTRL Packets Txed = %d",
        ophmr->stats.numQueryTxed + ophmr->stats.numReplySent
        ophmr->stats.numAckSent);
GLOMO_PrintStat(node, "RoutingOphmr", buf);
sprintf(buf, "Number of Data Txed = %d",
        ophmr->stats.numDataTxed);
GLOMO_PrintStat(node, "RoutingOphmr", buf);
sprintf(buf, "Number of Data Packets Originated = %d",
        ophmr->stats.numDataSent);
GLOMO_PrintStat(node, "RoutingOphmr", buf);

sprintf(buf, "Number of Data Packets Supposed to be Received = %d",
        ophmr->stats.numDataToReceive);
GLOMO_PrintStat(node, "RoutingOphmr", buf);

sprintf(buf, "Number of Data Packets Received = %d",
        ophmr->stats.numDataReceived);
GLOMO_PrintStat(node, "RoutingOphmr", buf);
}

```

```

void RoutingOphmrHandleData(GlomoNode *node, Message *msg)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;
    IpHeaderType *ipHdr = (IpHeaderType
*)GLOMO_MsgReturnPacket(msg);
    NODE_ADDR sourceAddress;
    NODE_ADDR destinationAddress;
    unsigned char IpProtocol;
    unsigned int ttl;
    NetworkQueueingPriorityType priority;

```

```

clocktype delay;
NODE_ADDR srcAddr = ipHdr->ip_src;
NODE_ADDR mcastAddr = ipHdr->ip_dst;
OphmrIpOptionType option = GetOphmrIpOptionField(msg);
Message *newMsg = NULL;

if (option.query == TRUE)
{
    RoutingOphmrHandleJoinQuery(node, msg);
    return;
}

if (!RoutingOphmrLookupMessageCache(srcAddr, option.seqNumber,
    &ophmr->messageCache))
{
    RoutingOphmrInsertMessageCache(node, srcAddr, option.seqNumber,
    &ophmr->messageCache);

    if (RoutingOphmrLookupMembership(mcastAddr, &ophmr-
>memberFlag))
    {
        #ifdef DEBUG
        printf("Node %ld received DATA\n", node->nodeAddr);
        #endif
        #ifdef DEBUG
        printf("  Member got it!\n");
        #endif
        ophmr->stats.numDataReceived++;
        newMsg = GLOMO_MsgCopy(node, msg);
        NetworkIpRemoveIpHeader(node, newMsg, &sourceAddress,
            &destinationAddress, &priority, &IpProtocol, &ttl);
        SendToUdp(node, newMsg, priority, sourceAddress,
            destinationAddress);
    }
    if (RoutingOphmrLookupFgFlag(mcastAddr, &ophmr->fgFlag))
    {
        #ifdef DEBUG
        printf("Node %ld received DATA\n", node->nodeAddr);
        #endif
        #ifdef DEBUG

```

```

printf("  FG. Forwarding it\n");
#endif
option.lastAddr = node->nodeAddr;
option.hopCount++;
SetOphmrIpOptionField(msg, &option);
delay = pc_erand(node->seed) * OPHMR_BROADCAST_JITTER;
NetworkIpSendPacketToMacLayerWithDelay(
    node, msg, DEFAULT_INTERFACE, ANY_DEST, delay);

    ophmr->stats.numDataTxed++;
}
}

else
{
    GLOMO_MsgFree(node, msg);
}
}

void RoutingOphmrHandleJoinQuery(GlomoNode *node, Message *msg)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;

    IpHeaderType *ipHdr = (IpHeaderType
*)GLOMO_MsgReturnPacket(msg);
    NODE_ADDR sourceAddress;
    NODE_ADDR destinationAddress;
    unsigned char IpProtocol;
    unsigned int ttl;
    NetworkQueueingPriorityType priority;
    clocktype delay, jrDelay;
    OPHMR_MT_Node *mcastEntry;
    OPHMR_RPT_Node *mEntry;
    NODE_ADDR srcAddr = ipHdr->ip_src;
    NODE_ADDR mcastAddr = ipHdr->ip_dst;
    OphmrIpOptionType option = GetOphmrIpOptionField(msg);
    Message *newMsg = NULL;
    if (!RoutingOphmrLookupMessageCache(
        srcAddr, option.seqNumber, &ophmr->messageCache))

```

```

{
    #ifdef DEBUG
    printf("Node %ld received Join Query from %d\n", node->nodeAddr,
        option.lastAddr);
    #endif
    RoutingOphmrInsertMessageCache(
        node, srcAddr, option.seqNumber, &ophmr->messageCache);
    {
        #ifdef DEBUG
        printf(" Member got it!\n");
        #endif
        ophmr->stats.numDataReceived++;
        if (option.hopCount < OPHMR_MAX_HOP)
        {
            option.lastAddr = node->nodeAddr;
            option.hopCount++;
            SetOphmrIpOptionField(msg, &option);
            delay = pc_erand(node->seed) *
                OPHMR_BROADCAST_JITTER;
            NetworkIpSendPacketToMacLayerWithDelay(
                node, msg, DEFAULT_INTERFACE, ANY_DEST, delay);

            ophmr->stats.numDataTxed++;
            ophmr->stats.numQueryTxed++;

            #ifdef DEBUG
            printf(" Relaying it\n");
            #endif

        }
        newMsg = GLOMO_MsgCopy(node, msg);
        NetworkIpRemoveIpHeader(node, newMsg, &sourceAddress,
            &destinationAddress, &priority, &IpProtocol, &ttl);

        SendToUdp(node, newMsg, priority, sourceAddress,
            destinationAddress);

        mcastEntry = RoutingOphmrGetMTEntry(mcastAddr, &ophmr-
            >memberTable);
    }
}

```

```

RoutingOphmrCheckSourceExpired(mcastAddr, &ophmr-
    >memberTable);

mEntry = RoutingOphmrGetRPTEntry(mcastAddr, &ophmr-
    >responseTable);

jrDelay = pc_erand(node->seed) * OPHMR_JR_JITTER;

RoutingOphmrSetTimer(
    node, MSG_NETWORK_SendReply, mcastAddr, jrDelay);
}
else if (option.hopCount < OPHMR_MAX_HOP)
{
    option.lastAddr = node->nodeAddr;
    option.hopCount++;
    SetOphmrIpOptionField(msg, &option);

    delay = pc_erand(node->seed) * OPHMR_BROADCAST_JITTER;

    NetworkIpSendPacketToMacLayerWithDelay(
        node, msg, DEFAULT_INTERFACE, ANY_DEST, delay);

    ophmr->stats.numDataTxed++;
    ophmr->stats.numQueryTxed++;
    #ifdef DEBUG
    printf("  Relaying it\n");
    #endif

}
else
{
    GLOMO_MsgFree(node, msg);
}

else
{
    GLOMO_MsgFree(node, msg);
}

```

```

void RoutingOphmrHandleReply(GlomoNode *node, Message *msg,
    NODE_ADDR lastAddr, NODE_ADDR mcastAddr)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;
    Message *newMsg;
    OPHMR_Ack *ackPkt;
    OPHMR_JoinReply *replyPkt = (OPHMR_JoinReply
        *)GLOMO_MsgReturnPacket(msg);
    OPHMR_TT_Node *mcastEntry;
    OPHMR_RPT_Node *mEntry;
    clocktype delay;
    int i;
    char *pktPtr;
    int pktSize = sizeof(OPHMR_Ack);
    BOOL changed = FALSE;

    for (i = 0; i < replyPkt->count; i++)
    {
        if (replyPkt->ackReq[i] && replyPkt->nextAddr[i] == node->nodeAddr)
        {
            newMsg = GLOMO_MsgAlloc(
                node, GLOMO_MAC_LAYER, 0, MSG_MAC_FromNetwork);
            GLOMO_MsgPacketAlloc(node, newMsg, pktSize);

            pktPtr = (char *) GLOMO_MsgReturnPacket(newMsg);
            ackPkt = (OPHMR_Ack *) pktPtr;
            ackPkt->pktType = OPHMR_ACK;
            ackPkt->mcastAddr = mcastAddr;
            ackPkt->srcAddr = replyPkt->srcAddr[i];
            NetworkIpSendRawGlomoMessageToMacLayer(node, newMsg,
                lastAddr, CONTROL, IPPROTO_OPHMR, 1,
                DEFAULT_INTERFACE, lastAddr);
            ophmr->stats.numAckSent++;
#ifdef DEBUG
            printf("Node %ld received a Join Reply from node %ld\n", node-
                >nodeAddr, lastAddr);
#endif
        }
    }
}

```

```

#ifdef DEBUG
printf("  sending Exp Ack to node %ld\n", lastAddr);
printf("  mcast = %u, src = %d\n", ackPkt->mcastAddr,      ackPkt-
      >srcAddr);
#endif
}
if (replyPkt->IAmFG)
{
  RoutingOphmrDeleteAckTable(mcastAddr,
      replyPkt->srcAddr[i], lastAddr, &ophmr->ackTable);
}
if (replyPkt->nextAddr[i] == node->nodeAddr &&
    replyPkt->srcAddr[i] != node->nodeAddr)
{
  #ifdef DEBUG
  printf("Node %ld received a Join Reply from node %ld\n", node-
      >nodeAddr, lastAddr);
  #endif
  #ifdef DEBUG
  printf("  I'm a FG!\n");
  #endif
  if (RoutingOphmrLookupFgFlag(mcastAddr, &ophmr->fgFlag))
  {
    RoutingOphmrUpdateFgFlag(mcastAddr, &ophmr->fgFlag);
  }
  else
  {
    RoutingOphmrSetFgFlag(mcastAddr, &ophmr->fgFlag);
  }

  mcastEntry = RoutingOphmrGetTTEEntry(
      mcastAddr, &ophmr->tempTable);
  if (mcastEntry == NULL)
  {
    #ifdef DEBUG
    #endif
    RoutingOphmrInsertTempTable(
        mcastAddr, replyPkt->srcAddr[i], &ophmr->tempTable);
  }
  else

```

```

{
    #ifndef DEBUG
    #endif
    RoutingOphmrInsertTempSource(replyPkt->srcAddr[i],
        mcastEntry);
}

if (RoutingOphmrCheckTempChanged(
    mcastAddr, &ophmr->tempTable))
{
    #ifndef DEBUG
    printf("    Temp changed!\n");
    #endif
    mEntry = RoutingOphmrGetRPTEntry(
        mcastAddr, &ophmr->responseTable);
    changed = TRUE;
}

RoutingOphmrCheckTempExpired(
    mcastAddr, &ophmr->tempTable);
}
}

if (changed && !RoutingOphmrCheckTempSent(
    mcastAddr, &ophmr->tempTable))
{
    #ifndef DEBUG
    printf("    Changed and temp sent!\n");
    #endif
    RoutingOphmrSetTempSent(mcastAddr, &ophmr->tempTable);
    delay = pc_erand(node->seed) * OPHMR_JR_JITTER +
        OPHMR_JR_PAUSE_TIME;
    RoutingOphmrSetTimer(
        node, MSG_NETWORK_SendReply, mcastAddr, delay);
}

if (RoutingOphmrLookupFgFlag(mcastAddr, &ophmr->fgFlag))
{
    RoutingOphmrSetTimer(

```

```

        node, MSG_NETWORK_CheckFg, mcastAddr,
OPHMR_FG_TIMEOUT);
    }

    GLOMO_MsgFree(node, msg);

}

void RoutingOphmrHandleAck(GlomoNode *node, Message *msg,
    NODE_ADDR lastAddr, NODE_ADDR targetAddr)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
        GlomoRoutingOphmr* ophmr =
        (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;

    OPHMR_Ack *ackPkt = (OPHMR_Ack
*)GLOMO_MsgReturnPacket(msg);

    if (targetAddr == node->nodeAddr)
    {
        #ifdef DEBUG
        printf("Node %ld got Ack from node %ld\n", node->nodeAddr,
            lastAddr);
        #endif
        RoutingOphmrDeleteAckTable(ackPkt->mcastAddr, ackPkt->srcAddr,
            lastAddr, &ophmr->ackTable);
    }

    GLOMO_MsgFree(node, msg);
}

void RoutingOphmrInitMessageCache(OPHMR_MC *messageCache)
{
    messageCache->front = NULL;
    messageCache->rear = NULL;
    messageCache->size = 0;
}

void RoutingOphmrInitSent(OPHMR_SS *sentTable)
{

```

```

sentTable->head = NULL;
sentTable->size = 0;
}
void RoutingOphmrInitStats(GlomoNode *node)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;
    ophmr->stats.numQueryTxed = 0;

    ophmr->stats.numReplySent = 0;

    ophmr->stats.numAckSent = 0;
    ophmr->stats.numDataSent = 0;
    ophmr->stats.numDataReceived = 0;
    ophmr->stats.numDataToReceive = 0;
    ophmr->stats.numDataTxed = 0;
}

void RoutingOphmrDeleteMsgCache(OPHMR_MC *messageCache)
{
    OPHMR_MC_Node *toFree;
    toFree = messageCache->front;
    messageCache->front = toFree->next;
    pc_free(toFree);
    --(messageCache->size);

    if (messageCache->size == 0)
    {
        messageCache->rear = NULL;
    }
}

BOOL RoutingOphmrLookupFgFlag(NODE_ADDR mcaddr,
OPHMR_FgFlag *fgFlag)
{
    OPHMR_FF_Node *current;
    if (fgFlag->size == 0)

```

```

{
    return (FALSE);
}

for (current = fgFlag->head;
     current != NULL; current = current->next)
{
    if (current->mcastAddr == mcastAddr)
    {
        return (TRUE);
    }
}
return (FALSE);
}

```

```

BOOL RoutingOphmrLookupSentTable(NODE_ADDR mcastAddr,
    OPHMR_SS *sentTable)

```

```

{
    OPHMR_SS_Node *current;

    if (sentTable->size == 0)
    {
        return (FALSE);
    }

    for (current = sentTable->head;
         current != NULL; current = current->next)
    {
        if (current->mcastAddr == mcastAddr)
        {
            return (TRUE);
        }
    }

    return (FALSE);
}

```

```

BOOL RoutingOphmrCheckFgExpired(NODE_ADDR mcastAddr,
    OPHMR_FgFlag *fgFlag)

```

```

OPHMR_FF_Node *current;

if (fgFlag->size == 0 || fgFlag->head == NULL)
{
    return (FALSE);
}

for (current = fgFlag->head;
     current != NULL; current = current->next)
{
    if (current->mcastAddr == mcastAddr &&
        simclock() - current->timestamp >= OPHMR_FG_TIMEOUT)
    {
        return (TRUE);
    }
}

return (FALSE);

void RoutingOphmrInsertAckTable(NODE_ADDR mcastAddr,
    OPHMR_AT *ackTable, OPHMR_JoinReply *reply)
{
    ++(ackTable->size);
    ackTable->head = RoutingOphmrInsertATInOrder(
        mcastAddr, ackTable->head, reply);
}

BOOL RoutingOphmrCheckASExist(NODE_ADDR srcAddr,
    OPHMR_AT_Node *mcast)
{
    OPHMR_AT_Snode *current;
    if (mcast->size == 0 || mcast->head == NULL)
    {
        return (FALSE);
    }
    for (current = mcast->head;
         current != NULL && current->srcAddr <= srcAddr;
         current = current->next)

```

```

{
    if (current->srcAddr == srcAddr)
    {
        return (TRUE);
    }
}
return (FALSE);
}

```

```

OPHMR_AT_Snode *RoutingOphmrInsertAckSource(
    NODE_ADDR srcAddr, NODE_ADDR nextAddr,
    OPHMR_AT_Node *mcast)
{
    OPHMR_AT_Snode *current;
    if (!RoutingOphmrCheckASExist(srcAddr, mcast))
    {
        ++(mcast->size);
        mcast->lastSent = simclock();
        mcast->head = RoutingOphmrInsertASInOrder(
            srcAddr, nextAddr, mcast->head, NULL);
        return (mcast->head);
    }
    else
    {
        for (current = mcast->head;
            current != NULL;
            current = current->next)
        {
            if (current->srcAddr == srcAddr)
            {
                current->nextAddr = nextAddr;
                current->numTx++;
                break;
            }
        }
        return (NULL);
    }
}
}

```

```

int RoutingOphmrGetAckCount(NODE_ADDR mcastAddr, OPHMR_AT
*ackTable)
{
    OPHMR_AT_Node *current;

    if (ackTable->size == 0 || ackTable->head == NULL)
    {
        return (0);
    }

    for (current = ackTable->head;
        current != NULL && current->mcastAddr <= mcastAddr;
        current = current->next)
    {
        if (current->mcastAddr == mcastAddr)
        {
            return (current->size);
        }
    }

    return (0);
}

```

```

void RoutingOphmrJoinGroup(GlomoNode *node, NODE_ADDR
mcastAddr)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *)
        node->networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;
    if (!RoutingOphmrLookupMembership(mcastAddr, &ophmr-
>memberFlag))
    {
        #ifdef DEBUG
        printf("Node %ld joined group %u\n", node->nodeAddr, mcastAddr);
        #endif
        RoutingOphmrSetMemberFlag(mcastAddr, &ophmr->memberFlag);
    }
}

```

```

void RoutingOphmrLeaveGroup(GlomoNode *node, NODE_ADDR
mcastAddr)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;
    if (RoutingOphmrLookupMembership(mcastAddr, &ophmr-
>memberFlag))
    {
        #ifdef DEBUG
        printf("Node %ld left group %u\n", node->nodeAddr, mcastAddr);
        #endif
        RoutingOphmrResetMemberFlag(mcastAddr, &ophmr->memberFlag);
        RoutingOphmrDeleteSourceSent(mcastAddr, &ophmr->sentTable);
        RoutingOphmrDeleteMemberTable(mcastAddr, &ophmr-
>memberTable);
    }
}

BOOL RoutingOphmrCheckAckTable(NODE_ADDR mcastAddr,
OPHMR_AT *ackTable)
{
    OPHMR_AT_Node *current;
    OPHMR_AT_Snode *curSrc;
    if (ackTable->size == 0 || ackTable->head == NULL)
    {
        return (FALSE);
    }
    for (current = ackTable->head; current != NULL; current = current->next)
    {
        if (current->mcastAddr == mcastAddr)
        {
            if (current->size == 0 || current->head == NULL)
            {
                return (FALSE);
            }
            for (curSrc = current->head;

```

```

    curSrc != NULL;
    curSrc = curSrc->next)
{
    if (curSrc->numTx < OPHMR_MAX_NUM_TX)
    {
        return (TRUE);
    }
}
}
return (FALSE);
}

```

```

void RoutingOphmrDeleteAckTable(NODE_ADDR mcastAddr,
NODE_ADDR srcAddr,
    NODE_ADDR lastAddr, OPHMR_AT *ackTable)
{
    OPHMR_AT_Node *current;
    OPHMR_AT_Snode *curSrc;
    OPHMR_AT_Snode *toFree;
    for (current = ackTable->head; current != NULL; current = current-
        >next)
    {
        if (current->mcastAddr == mcastAddr)
        {
            if (current->size == 0 || current->head == NULL)
            {
                return;
            }
            for (curSrc = current->head;
                curSrc != NULL;
                curSrc = curSrc->next)
            {
                if (curSrc->srcAddr == srcAddr &&
                    curSrc->nextAddr == lastAddr)
                {
                    toFree = curSrc;

                    if (curSrc->prev == NULL && curSrc->next == NULL)
                    {

```

```

    current->head = curSrc->next;
}
else if (curSrc->prev == NULL)
{
    curSrc->next->prev = curSrc->prev;
    current->head = curSrc->next;
}
else if (curSrc->next == NULL)
{
    curSrc->prev->next = curSrc->next;
}
else
{
    curSrc->prev->next = curSrc->next;
    curSrc->next->prev = curSrc->prev;
}

pc_free(toFree);
--(current->size);
}

else if (curSrc->numTx >= OPHMR_MAX_NUM_TX)
{
    toFree = curSrc;

    if (curSrc->prev == NULL && curSrc->next == NULL)
    {
        current->head = curSrc->next;
    }
    else if (curSrc->prev == NULL)
    {
        curSrc->next->prev = curSrc->prev;
        current->head = curSrc->next;
    }
    else if (curSrc->next == NULL)
    {
        curSrc->prev->next = curSrc->next;
    }
    else
    {

```

```

    curSrc->prev->next = curSrc->next;
    curSrc->next->prev = curSrc->prev;
}

```

```

    pc_free(toFree);
    --(current->size);
}
}
}
}
}
}

```

```

void RoutingOphmrInitAckTable(OPHMR_AT *ackTable)

```

```

{
    ackTable->head = NULL;
    ackTable->size = 0;
}

```

```

if (current->next != NULL && current->next->mcastAddr == mcastAddr)

```

```

{
    toFree = current->next;
    current->next = toFree->next;
    pc_free(toFree);
    --(tempTable->size);
}
}
}
}

```

```

void RoutingOphmrSetTimer(GlomoNode *node, long eventType,
    NODE_ADDR mcastAddr, clocktype delay)

```

```

{
    Message *msg;
    NODE_ADDR *info;
    msg = GLOMO_MsgAlloc(node, GLOMO_NETWORK_LAYER,
        ROUTING_PROTOCOL_OPHMR, eventType);
    GLOMO_MsgInfoAlloc(node, msg, sizeof(NODE_ADDR));
    info = (NODE_ADDR *) GLOMO_MsgReturnInfo(msg);
    *info = mcastAddr;
    GLOMO_MsgSend(node, msg, delay);
}

```

```
}
```

```
void RoutingOphmrSendReply(GlomoNode *node, NODE_ADDR  
mcastAddr,  
                          OPHMR_MT *memberTable, OPHMR_TT *tempTable)
```

```
{
```

```
  GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-  
                              >networkData.networkVar;  
  GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-  
                              >routingProtocol;
```

```
  Message *msg;
```

```
  OPHMR_JoinReply *replyPkt;
```

```
  OPHMR_AT_Node *mcastEntry;
```

```
  OPHMR_MT_Node *current;
```

```
  OPHMR_MT_Snode *curSrc;
```

```
  OPHMR_TT_Node *curNode;
```

```
  OPHMR_TT_Snode *curSrcNode;
```

```
  BOOL match;
```

```
  BOOL done;
```

```
  int i, j, k, l, m, n;
```

```
  char *pktPtr;
```

```
  int pktSize = sizeof(OPHMR_JoinReply);
```

```
  clocktype delay;
```

```
  BOOL needAck = FALSE;
```

```
  GLOMO_MsgPacketAlloc(node, msg, pktSize);
```

```
  pktPtr = (char *) GLOMO_MsgReturnPacket(msg);
```

```
  replyPkt = (OPHMR_JoinReply *) pktPtr;
```

```
  replyPkt->pktType = OPHMR_JOIN_REPLY;
```

```
  RoutingOphmrSetTempSent(mcastAddr, &ophmr->tempTable);
```

```
  if (RoutingOphmrLookupMemberTable(mcastAddr, &ophmr-  
      >memberTable) && RoutingOphmrCheckLastSent(mcastAddr,  
      &ophmr->memberTable))
```

```
  {
```

```
    #ifdef DEBUG
```

```
      printf("    There are members!\n");
```

```
    #endif
```

```
    replyPkt->count = RoutingOphmrGetMemberCount(  
                  mcastAddr, &ophmr->memberTable);
```

```
    #ifdef DEBUG
```

```
      #endif
```

```

replyPkt->IAmFG = RoutingOphmrLookupFgFlag(
    mcastAddr, &ophmr->fgFlag);
for (current = memberTable->head;
    current != NULL && current->mcastAddr <= mcastAddr;
    current = current->next)
{
    if (current->mcastAddr == mcastAddr)
    {
        for (i = 0, curSrc = current->head;
            i < current->size && curSrc != NULL;
            i++, curSrc = curSrc->next)
        {
            replyPkt->srcAddr[i] = curSrc->srcAddr;
            #ifdef DEBUG
            #endif
        }
    }
}

for (j = 0; j < replyPkt->count; j++)
{
    replyPkt->nextAddr[j] = RoutingOphmrGetNextNode(
        replyPkt->srcAddr[j], &ophmr->routeTable);
    #ifdef DEBUG
    #endif
    if (replyPkt->nextAddr[j] == replyPkt->srcAddr[j])
    {
        replyPkt->ackReq[j] = TRUE
        needAck = TRUE;
    }
    else
    {
        replyPkt->ackReq[j] = FALSE;
        needAck = TRUE;
    }
}

if (tempTable->size == 0 || tempTable->head == NULL)
{
}

```

```

else
{
    for (curNode = tempTable->head;
        curNode != NULL && curNode->mcastAddr <=
            mcastAddr; curNode = curNode->next)
    {
        if (curNode->mcastAddr == mcastAddr)
        {
            if (curNode->size == 0 || curNode->head == NULL)
            {
            }
            else
            {
                for (curSrcNode = curNode->head;
                    curSrcNode != NULL;
                    curSrcNode = curSrcNode->next)
                {
                    done = FALSE;
                    for (k = 0; k < replyPkt->count; k++)
                    {
                        if (replyPkt->srcAddr[k] ==
                            curSrcNode->srcAddr)
                        {
                            done = TRUE;
                            break;
                        }
                    }
                }

                if (done == FALSE)
                {
                    replyPkt->srcAddr[replyPkt->count] =
                        curSrcNode->srcAddr;
                }
                else
                {
                    replyPkt->ackReq[replyPkt->count] =

                                FALSE;
                }
                replyPkt->count++;
            }
        }
    }
}

```

```
    }
  }
}
```

```
#ifndef DEBUG
#endif
while (l < replyPkt->count)
{
  match = FALSE;
  match = RoutingOphmrCheckResponseMatch(
mcastAddr, replyPkt->srcAddr[l], &ophmr->responseTable);
  if (!match)
  {
    for (m = l; m < replyPkt->count; m++)
    {
      replyPkt->srcAddr[m] = replyPkt->srcAddr[m+1];
      replyPkt->nextAddr[m] = replyPkt->nextAddr[m+1];
      replyPkt->ackReq[m] = replyPkt->ackReq[m+1];
    }

    replyPkt->count--;
  }
  if (match)
  {
    l++;
  }
}
```

```
#ifndef DEBUG
#endif
if (replyPkt->count > 0)
{
  RoutingOphmrDeleteResponseTable(
  mcastAddr, &ophmr->responseTable);
```

```
#ifndef DEBUG
for (l = 0; l < replyPkt->count; l++)
```



```

RoutingOphmrSetTimer(node, MSG_NETWORK_SendReply,
                    mcastAddr, OPHMR_JR_PAUSE_TIME);
}

else if (RoutingOphmrLookupFgFlag(mcastAddr, &ophmr->fgFlag))
{
#ifdef DEBUG
printf("FG node %ld sending a Join Reply\n", node->nodeAddr);
#endif
replyPkt->count = RoutingOphmrGetTempCount(
    mcastAddr, &ophmr->tempTable);
replyPkt->IAmFG =
    RoutingOphmrLookupFgFlag(mcastAddr, &ophmr->fgFlag);
for (curNode = tempTable->head;
    curNode != NULL && curNode->mcastAddr <= mcastAddr;
    curNode = curNode->next)
{
    if (curNode->mcastAddr == mcastAddr)
    {
        for (i = 0, curSrcNode = curNode->head;
            i < curNode->size && curSrcNode != NULL;
            i++, curSrcNode = curSrcNode->next)
        {
            replyPkt->srcAddr[i] = curSrcNode->srcAddr;
        }
    }
}
for (j = 0; j < replyPkt->count; j++)
{
    replyPkt->nextAddr[j] = RoutingOphmrGetNextNode(
        replyPkt->srcAddr[j], &ophmr->routeTable);
    if (replyPkt->nextAddr[j] == replyPkt->srcAddr[j])
    {
        replyPkt->ackReq[j] = TRUE needAck = TRUE;
    }
    else
    {
        replyPkt->ackReq[j] = FALSE;
        needAck = TRUE;
    }
}
}

```

```

    }
}
k = 0;
while (k < replyPkt->count)
{
    match = FALSE;
    match = RoutingOphmrCheckResponseMatch(
mcastAddr, replyPkt->srcAddr[k], &ophmr->responseTable);
    if (!match)
    {
        for (l = k; l < replyPkt->count; l++)
        {
            replyPkt->srcAddr[l] = replyPkt->srcAddr[l+1];
            replyPkt->nextAddr[l] = replyPkt->nextAddr[l+1];
            replyPkt->ackReq[l] = replyPkt->ackReq[l+1];
        }
        replyPkt->count--;
    }
    if (match)
    {
        k++;
    }
}
NetworkIpSendRawGlomoMessageToMacLayer(
node, msg, mcastAddr, CONTROL, IPPROTO_OPHMR, 1,
DEFAULT_INTERFACE, ANY_DEST);
ophmr->stats.numReplySent++;
if (needAck)
{
    mcastEntry = RoutingOphmrGetATEntry(mcastAddr, &ophmr-
>ackTable);
    if (mcastEntry == NULL)
    {
        RoutingOphmrInsertAckTable(
            mcastAddr, &ophmr->ackTable, replyPkt);
    }
    else
    {
        for (m = 0; m < replyPkt->count; m++)
        {

```

```

RoutingOphmrInsertAckSource(
    replyPkt->srcAddr[m], replyPkt->nextAddr[m], mcastEntry);
}
}

```

```

RoutingOphmrSetTimer(
    node, MSG_NETWORK_CheckAked, mcastAddr,
    OPHMR_CHECKACK_INTERVAL);
}
}
}
}
}

```

```

void RoutingOphmrRetxReply(
    GlomoNode *node, NODE_ADDR mcastAddr, OPHMR_AT *ackTable)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;

    Message *msg;
    OPHMR_JoinReply *replyPkt;
    OPHMR_AT_Node *mcastEntry;
    OPHMR_AT_Node *current;
    OPHMR_AT_Snode *curSrc;
    int i, j;
    char *pktPtr;
    int pktSize = sizeof(OPHMR_JoinReply);
    clocktype delay;

    msg = GLOMO_MsgAlloc(node, GLOMO_MAC_LAYER, 0,
        MSG_MAC_FromNetwork);
    GLOMO_MsgPacketAlloc(node, msg, pktSize);
    pktPtr = (char *) GLOMO_MsgReturnPacket(msg);
    replyPkt = (OPHMR_JoinReply *) pktPtr;
    replyPkt->pktType = OPHMR_JOIN_REPLY;
    replyPkt->count = RoutingOphmrGetAckCount(mcastAddr, &ophmr-
        >ackTable);
    replyPkt->IAmFG = RoutingOphmrLookupFgFlag(mcastAddr, &ophmr-
        >fgFlag);
}

```

```

for (current = ackTable->head;
    current != NULL && current->mcastAddr <= mcastAddr;
    current = current->next)
{
    if (current->mcastAddr == mcastAddr)
    {
        for (i = 0, curSrc = current->head;
            i < current->size && curSrc != NULL;
            i++, curSrc = curSrc->next)
        {
            replyPkt->srcAddr[i] = curSrc->srcAddr;
            replyPkt->nextAddr[i] = curSrc->nextAddr;
            if (curSrc->numTx > 0 || curSrc->srcAddr == curSrc->nextAddr)
            {
                replyPkt->ackReq[i] = TRUE;
            }
            else
            {
                replyPkt->ackReq[i] = FALSE;
            }
        }
#ifdef DEBUG
        printf("src[%d] = %d, next[%d] = %d, ack[%d] = %d\n", i,
            replyPkt->srcAddr[i], i, replyPkt->nextAddr[i], i, replyPkt-
                >ackReq[i]);
#endif
    }
}
}
if (replyPkt->count > 0)
{
    delay = pc_erand(node->seed) * OPHMR_JR_RETX_JITTER;
    NetworkIpSendRawGlomoMessageToMacLayerWithDelay(
        node, msg, mcastAddr, CONTROL, IPPROTO_OPHMR, 1,
        DEFAULT_INTERFACE, ANY_DEST, delay);
    ophmr->stats.numReplySent++;
#ifdef DEBUG
    printf("Node %ld retransmitting a Join Reply\n", node->nodeAddr);
#endif
    RoutingOphmrSetTimer(

```

```

node, MSG_NETWORK_CheckAcked, mcastAddr,
    OPHMR_CHECKACK_INTERVAL);
mcastEntry = RoutingOphmrGetATEntry(mcastAddr, &ophmr-
    >ackTable);
if (mcastEntry == NULL)
{
    RoutingOphmrInsertAckTable(mcastAddr, &ophmr->ackTable,
        replyPkt);
}
else
{
    for (j = 0; j < replyPkt->count; j++)
    {
        RoutingOphmrInsertAckSource(replyPkt->srcAddr[j], replyPkt-
            >nextAddr[j], mcastEntry);
    }
}
}

}

void RoutingOphmrSendData(GlomoNode *node, Message *msg,
NODE_ADDR mcastAddr)
{
    GlomoNetworkIp* ipLayer = (GlomoNetworkIp *) node-
        >networkData.networkVar;
    GlomoRoutingOphmr* ophmr = (GlomoRoutingOphmr *) ipLayer-
        >routingProtocol;

    OphmrIpOptionType option;
    GLOMO_MsgSetLayer(msg, GLOMO_MAC_LAYER, 0);
    GLOMO_MsgSetEvent(msg, MSG_MAC_FromNetwork);
    AddCustomOphmrIpOptionFields(node, msg);
    option.query = FALSE;
    option.lastAddr = node->nodeAddr;
    option.seqNumber = RoutingOphmrGetSeq(node);
    option.hopCount = 1;
    NetworkIpSendPacketToMacLayer(node, msg,DEFAULT_INTERFACE,
    ANY_DEST);
    void RoutingOphmrSendQuery(GlomoNode *node, Message *msg,
    NODE_ADDR mcastAddr){

```



## 8.REFERENCES

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